



A Solution Approach for AS Relationships-aware Overlay Routing

<draft-asai-cross-domain-overlay-00>

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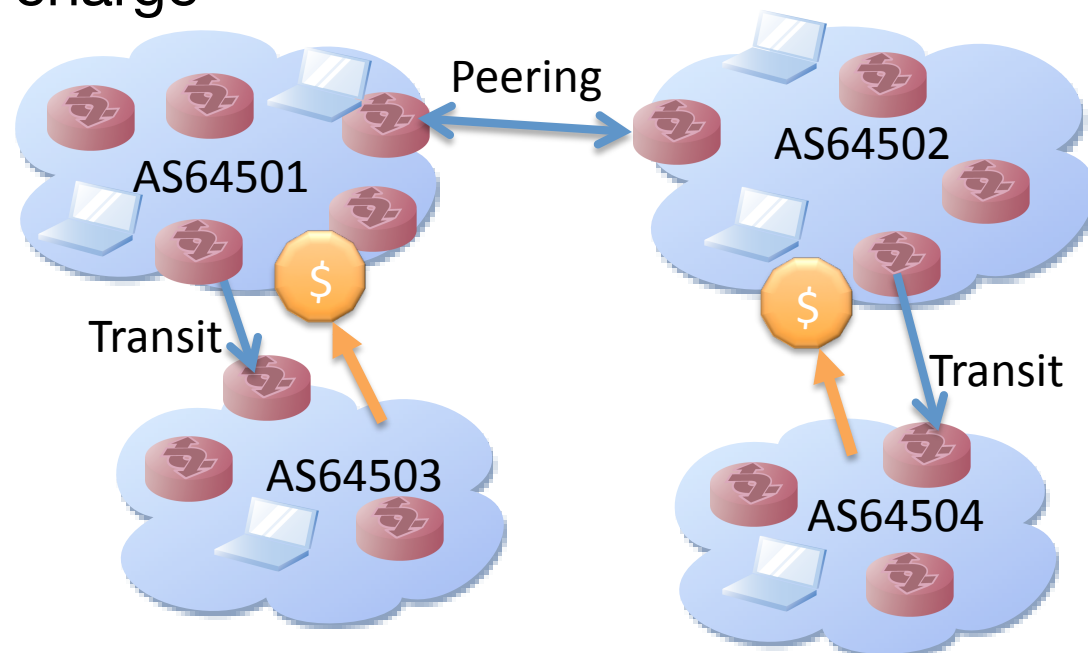
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The Internet

- Autonomous systems (ASes)
 - e.g., ISPs, companies, and universities
- Inter-AS economics
 - transit charge



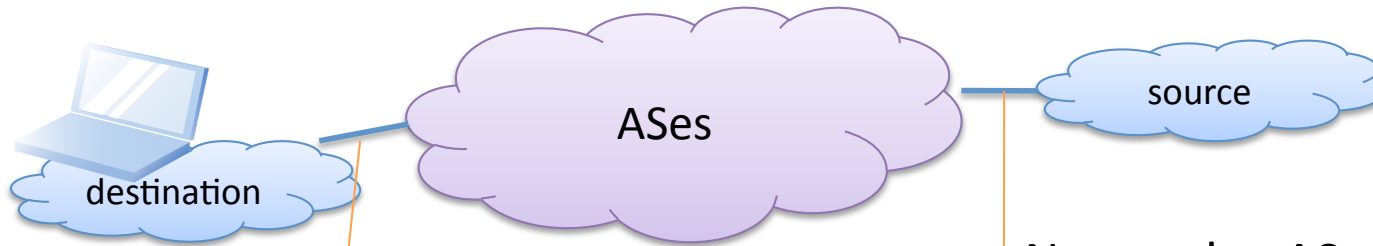
Motivation

- Reduction in transit traffic/charge
 - Why?
 - Most of ISPs providing their network to P2P nodes are residential ones.
 - i.e., not tier-1, but customers
 - Transit traffic costs more compared to peering or intra-domain traffic.

Approach

- Be aware of commercial relationships between ASes
 - in overlay networks
 - Similar to ALTO
 - but
 - focusing mainly on cross-domain traffic
 - with hiding ISP's confidential information as much as possible (i.e., minimum ISP cooperation)

Detail with valley-free path model



Note; edge ASes accommodate CDN peers.

		Source Edge		
		p2c	p2p	c2p
Destination Edge	p2c	⊘ +, +	⊘ +, 0	+ , -
	p2p	⊘ 0, +	0, 0	0, -
	c2p	- , +	- , 0	- , -

higher preference ←

p2c: provider to customer
c2p: customer to provider
p2p: peer to peer

the worst path: both ASes pay transit charge.



Requirements (1/2): AS relationships

- Information on AS relationships
 - Inference methods (N.B., assuming ISPs do not want to reveal AS relationships)
 - Path analysis [Gao2001, etc]
 - Adjacency analysis [Asai2010]
 - A method for provisioning this information to peers
 - » Inferred then provisioned from server(s)
 - *written in the draft*



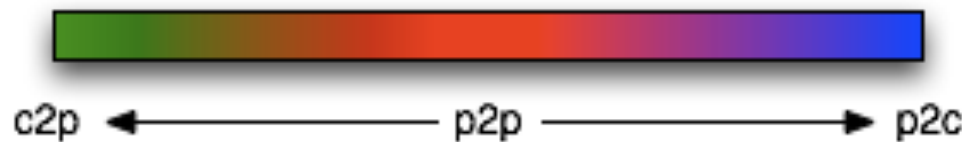
Requirements (2/2): Delivery cost for a certain path

- Delivery cost computation
 - End-to-end path (AS path)
 - A method for provisioning this information to peers
 - traceroute by peers
 - Provisioned by servers in ASes
 - » **written in the draft**
 - Function
 - i.e., $f(P) = (\text{transit charge on residential ASes})$
 - f : function, P : AS path

AS relationships inference

- Heuristics: Common approach
 - Degree (i.e., # of neighbors)
 - High degree \Leftrightarrow large
 - tends to be provider
 - Low degree \Leftrightarrow small
 - tends to be customer

Diff. in degree

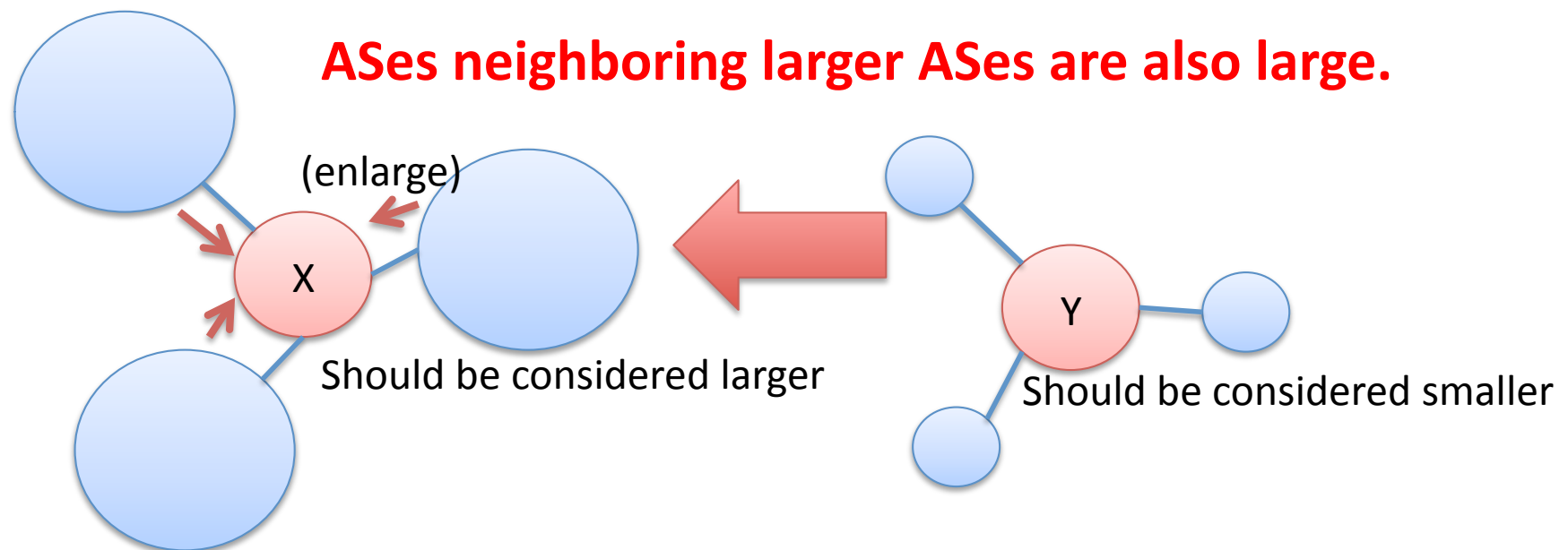


p2c: provider to customer
c2p: customer to provider
p2p: peer to peer

AS relationships inference (cont'd): Improving degree-based approach

Take into account the size of n -hop neighbor ASes

→ “Magnitude”

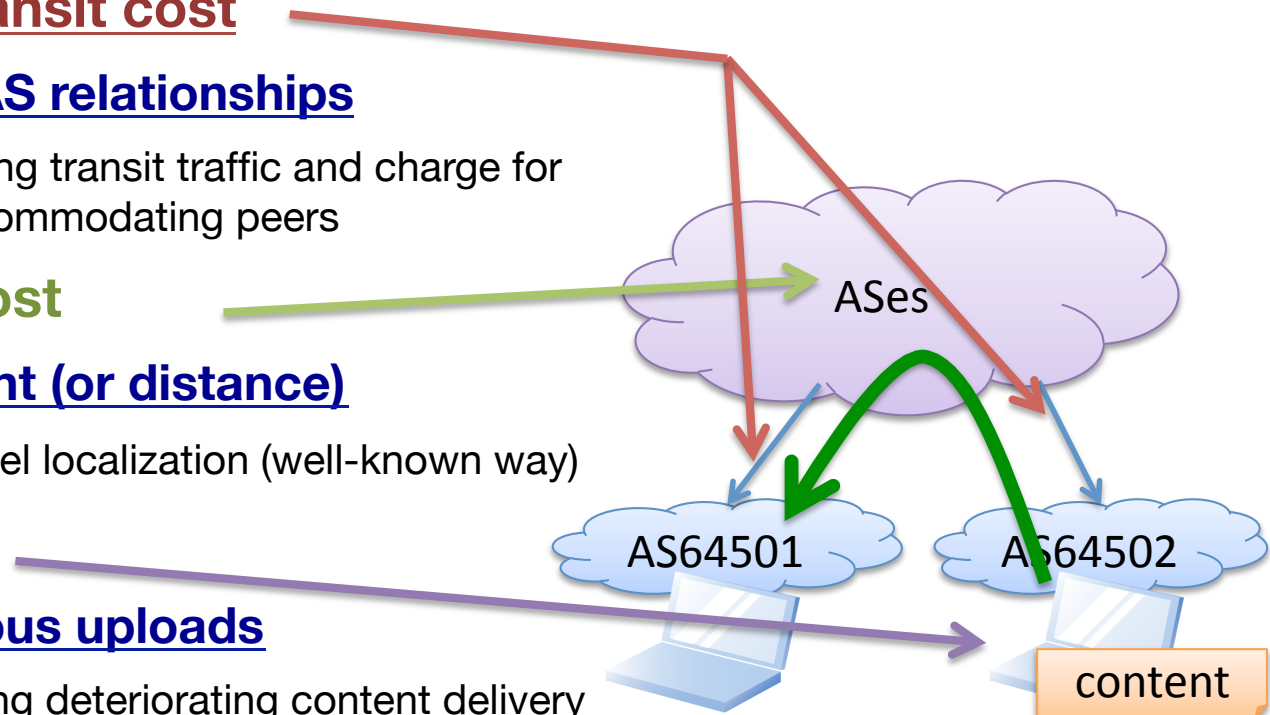


c.f. H. Asai et al., “Estimating AS Relationships for Application-Layer Traffic Optimization,” 2010

7.28, 2010 H. Asai, H. Esaki and T. Momose, “AS Relationships-aware Overlay Routing”

Example of cost function (e.g., in CDN)

- Metrics for peer selection
 - **(1) Inter-AS transit cost**
 - **Estimated AS relationships**
 - for reducing transit traffic and charge for ASes accommodating peers
 - **(2) Network cost**
 - **AS hop count (or distance)**
 - for AS-level localization (well-known way)
 - **(3) Quality**
 - **#simultaneous uploads**
 - for avoiding deteriorating content delivery throughput (very naïve parameter...)



Simulation result: High-cost transit traffic reduction

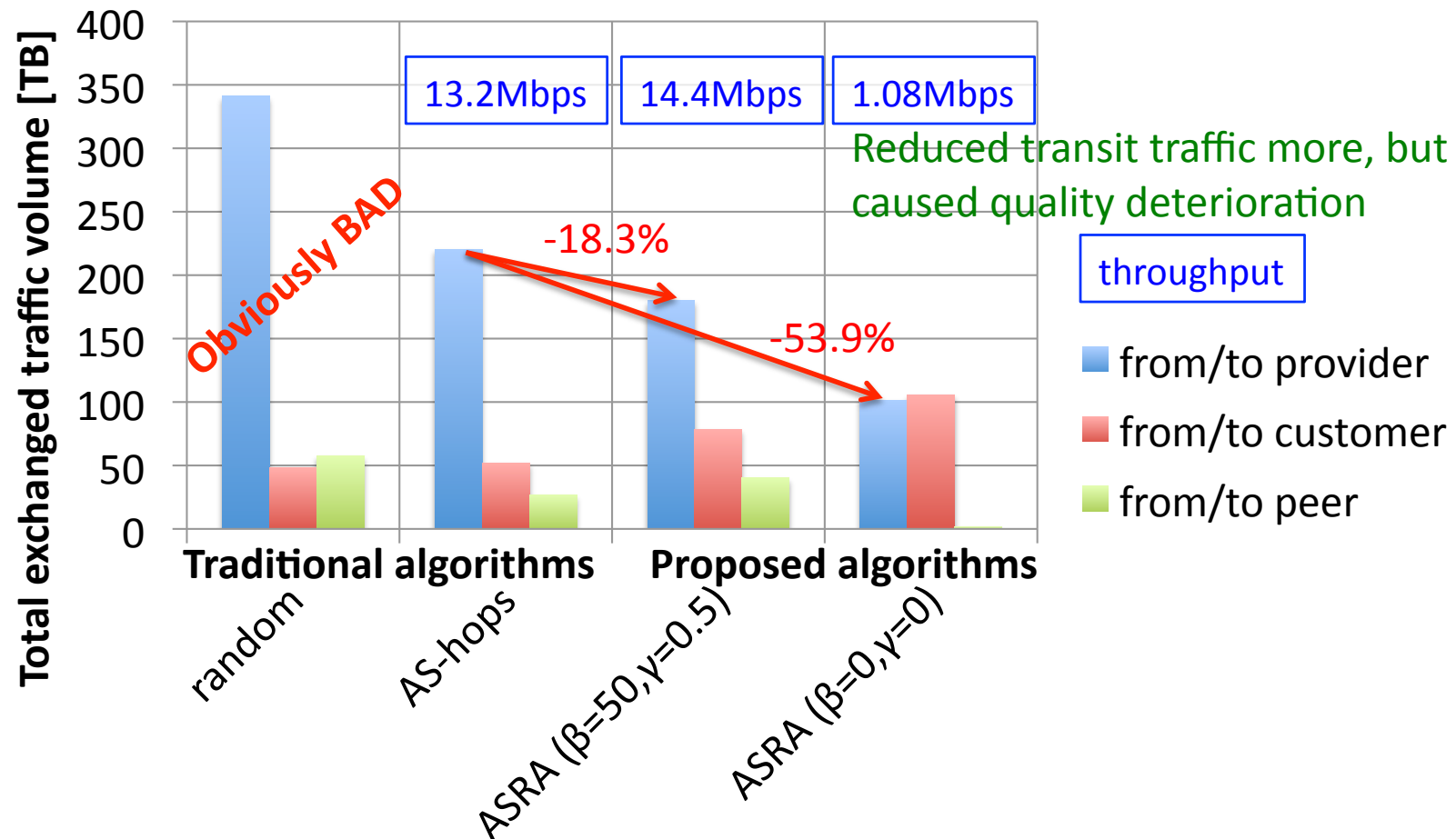


Fig. Breakdown of inter-domain traffic on ASes accommodating CDN peers



Simulation result: Transit charge reduction

N.B., provider ASes charges customer ASes for transit traffic based on the exchanged traffic volume.

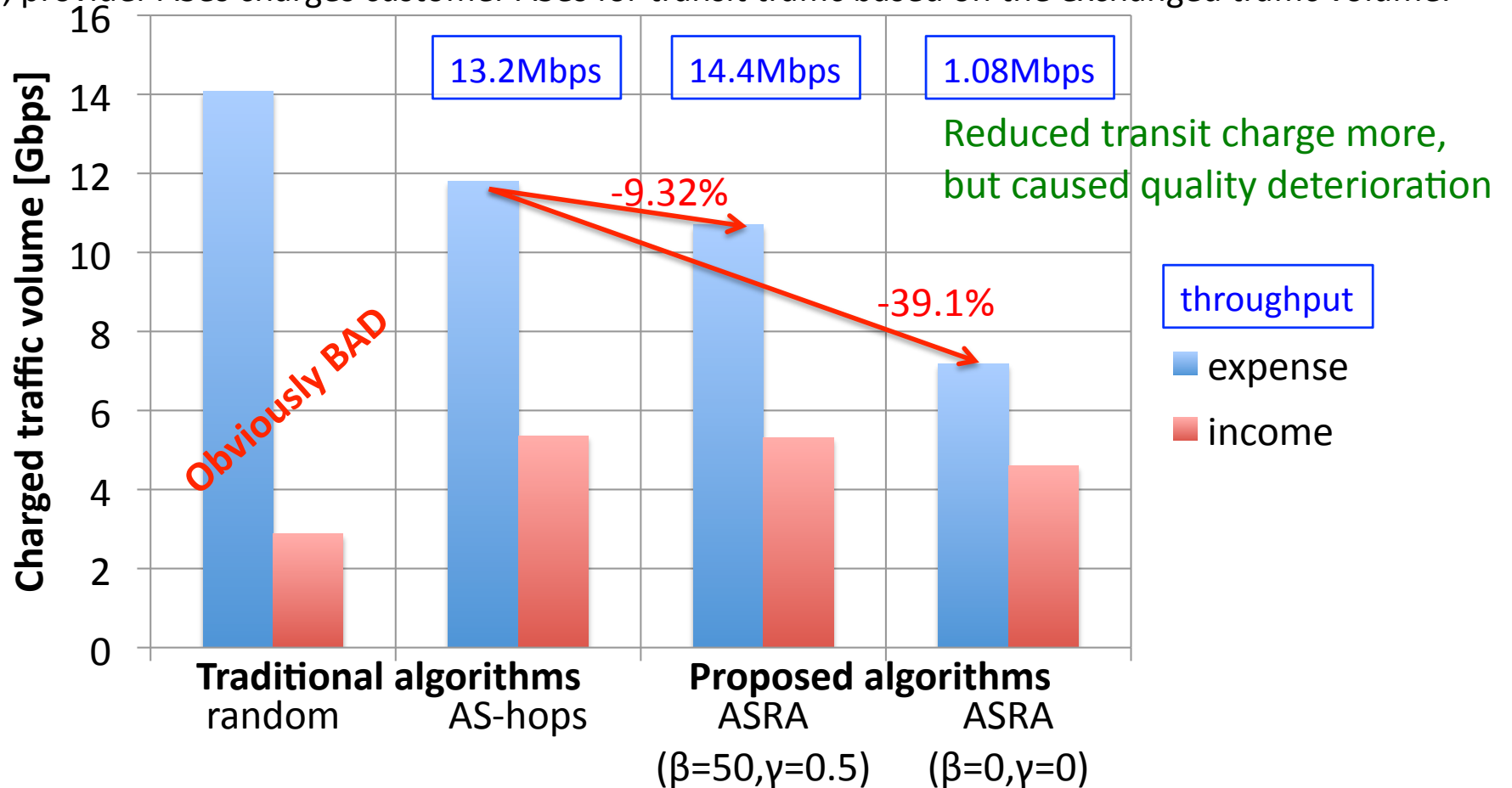
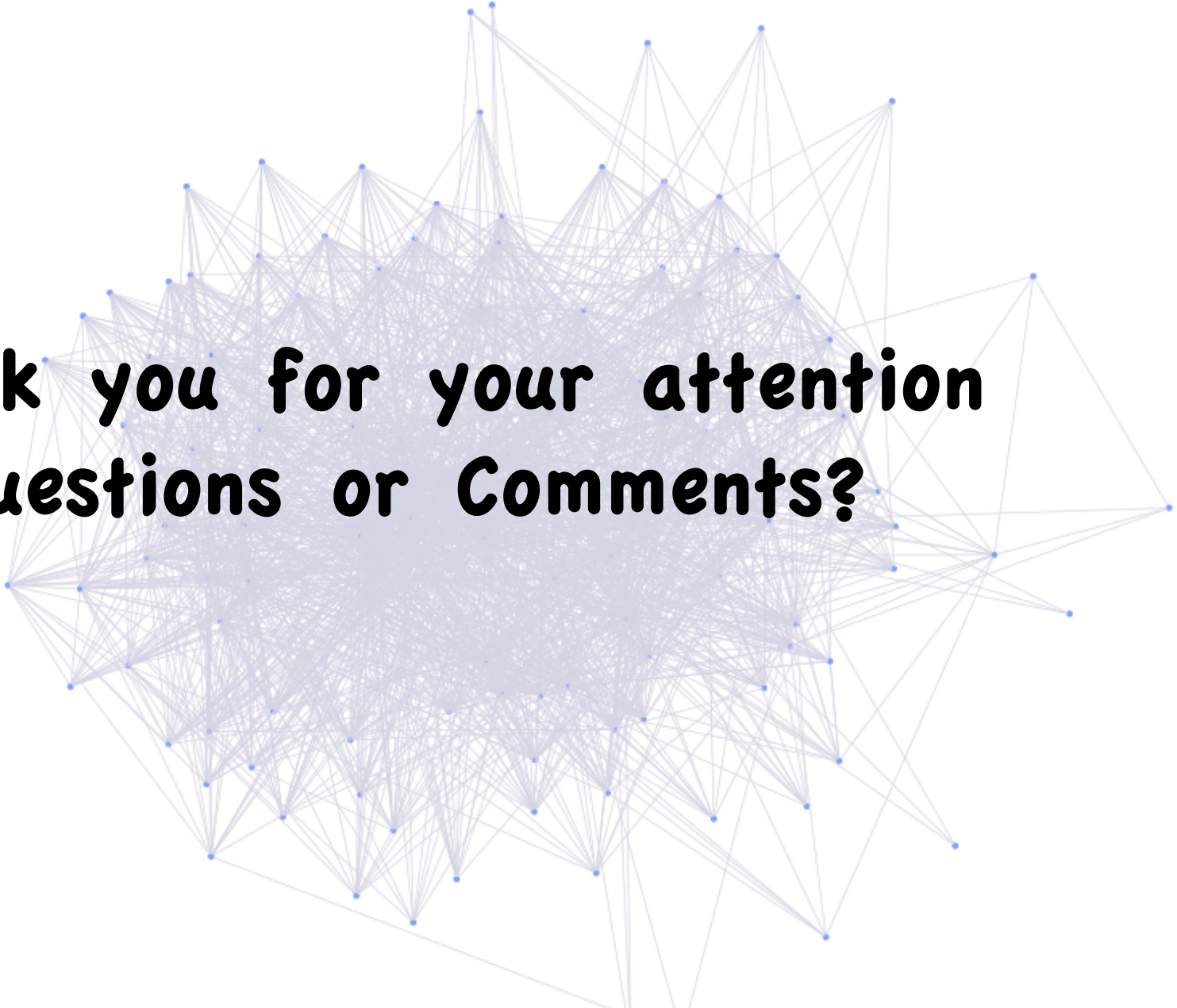


Fig. Charged transit traffic volume with the 95th percentile charging policy

Conclusion

- Approach for reducing transit traffic
 - Pros.
 - Not require ISPs' information (available by end-to-end)
 - Cons.
 - Efficiency depends on inference methods.
- TODO
 - Field experiment to evaluate the system
 - Discussion on deployment possibility and next step
 - Can ISPs provide AS paths?
 - Yes → How to deploy
 - No → traceroute-based approach or other alternatives?



**Thank you for your attention
Questions or Comments?**

BACKUP SLIDES: SIMULATION SETUP

Improved AS relationships inference

(1) Define a weighted AS adjacency matrix

$${}^n A := ({}^n a_{v_i v_j}) \quad \begin{array}{l} \text{(i) } n = 0 \\ \text{(ii) } n \geq 1, n \in \mathbb{Z} \end{array}$$

$${}^n a_{v_i v_j} = \begin{cases} 1 & : \text{ if AS } v_i \text{ and AS } v_j \text{ are adjacent} \\ 0 & : \text{ otherwise} \end{cases}$$

$${}^n a_{v_i v_j} = \begin{cases} (n-1)\rho_{v_j} & : \text{ if AS } v_i \text{ and AS } v_j \text{ are adjacent} \\ 0 & : \text{ otherwise} \end{cases}$$

(2) Convert the weighted AS adjacency matrix to a traffic transition matrix

$${}^n T := \left(\frac{{}^n a_{v_i v_j}}{\sum_{v_k} {}^n a_{v_i v_k}} \right)$$

(3) Calculate the left eigenvector of the traffic transition matrix corresponding to the maximum eigenvalue

$$\boxed{{}^n \boldsymbol{\rho}} = [{}^n \rho_{v_1}, \dots, {}^n \rho_{v_m}]^t : \text{ the left eigenvector}$$

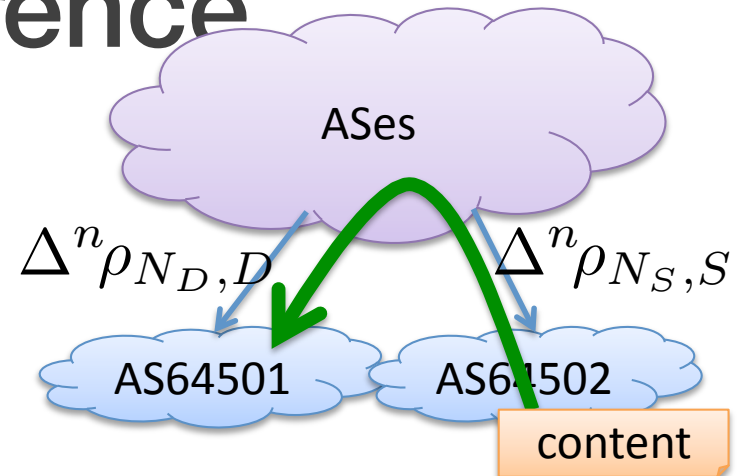
$$(\|{}^n \boldsymbol{\rho}\| = 1)$$

Peer selection preference

Peer Selection Preference

$$p(P) := \alpha c_i(P) + \beta c_t(P) + \gamma q(P)$$

α, β, γ : parameters



(1) Internal CDN cost; i.e., AS relationships

$$c_i(P) := \begin{cases} \epsilon \Delta^n \rho_{N_D, D} + (1 - \epsilon) \Delta^n \rho_{N_S, S} + 2 |\max^n \rho| & (S \neq D) \\ 0 & (S = D) \end{cases}$$

ϵ : weighting factor, $\epsilon = 0.5$

given (by traceroute etc.)

(2) Total network cost

$$c_t(P) := H_{S, D}$$

(3) Quality

$$q(P) := b_s^2$$

P : path from peer s in AS S to peer d in AS D
 N_S : neighbor (next hop) AS of AS S
 N_D : neighbor (previous hop) AS of AS D
 $H_{S, D}$: AS hop count from AS S to AS D
 b_s : the number of simultaneous uploads on peer s

Simulation setup

- Evaluation model
 - Request pattern
 - based on measured peer distribution in *BitTorrent*
 - Internet topology
 - CAIDA AS Relationships Dataset (10/08/2009)
- Algorithms
 - random: select one uniformly at random
 - AS-hops: minimize AS hop count
 - proposed algorithm ($\alpha = 1, \varepsilon = 0.5$): minimize preference
 - $(\beta, \gamma) = \{(0, 0), (0, 0.5), (50, 0), (50, 0.5)\}$

