

FRTR: A Scalable Mechanism for Global Routing Consistency

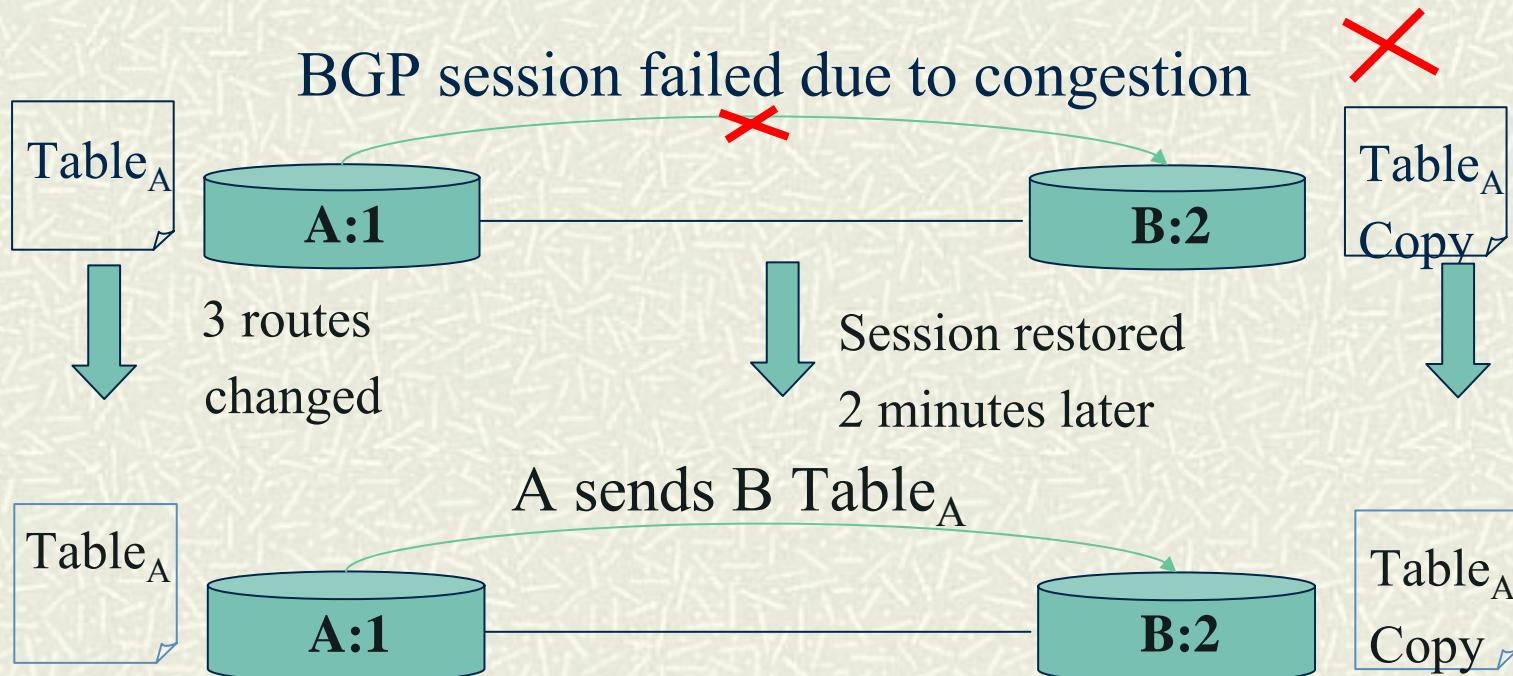
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Why does BGP need fast routing recovery?

BGP recovers inefficiently from transient session resets.



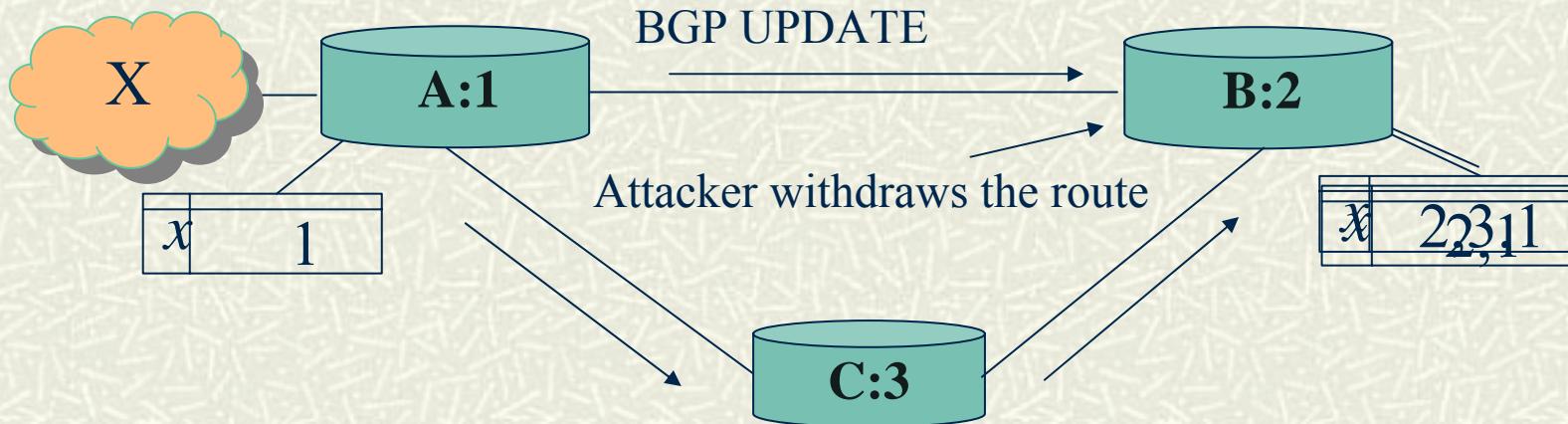
3 routes were changed, but A sent 120K+ routes (~5MB).

Why does BGP need fast routing recovery? (cont'd)

Faults and attacks can corrupt a routing table.

Originate route to network X

Route to X is through A



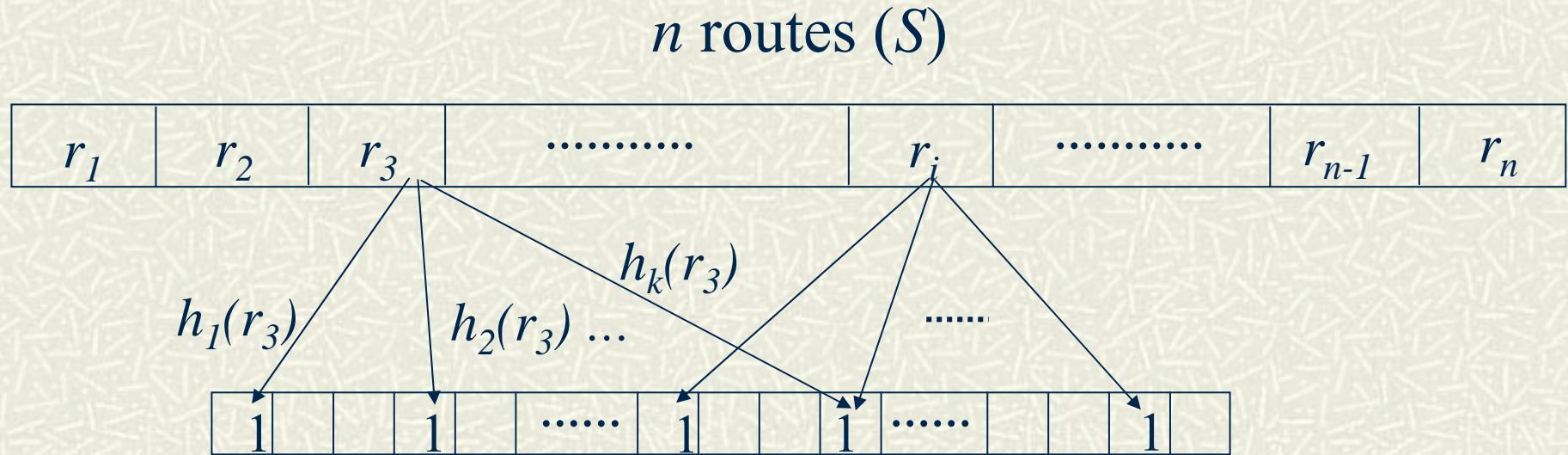
Attacker has successfully changed B's route to X!

Other fault sources:
hardware failures, human errors, software bugs

FRTR (Fast Routing Table Recovery)

- Use small Bloom filter to encode routing table;
- Exchange digests periodically.

Bloom Filter Digest



l -bit digest (d)

l/n : encoding ratio, a small constant, e.g. 5

$h_1(), h_2(), \dots, h_k()$: hash functions

Is r a member of S ? \Rightarrow check $d(h_j(r)) = 1?$ ($j = 1, 2, \dots, k$)

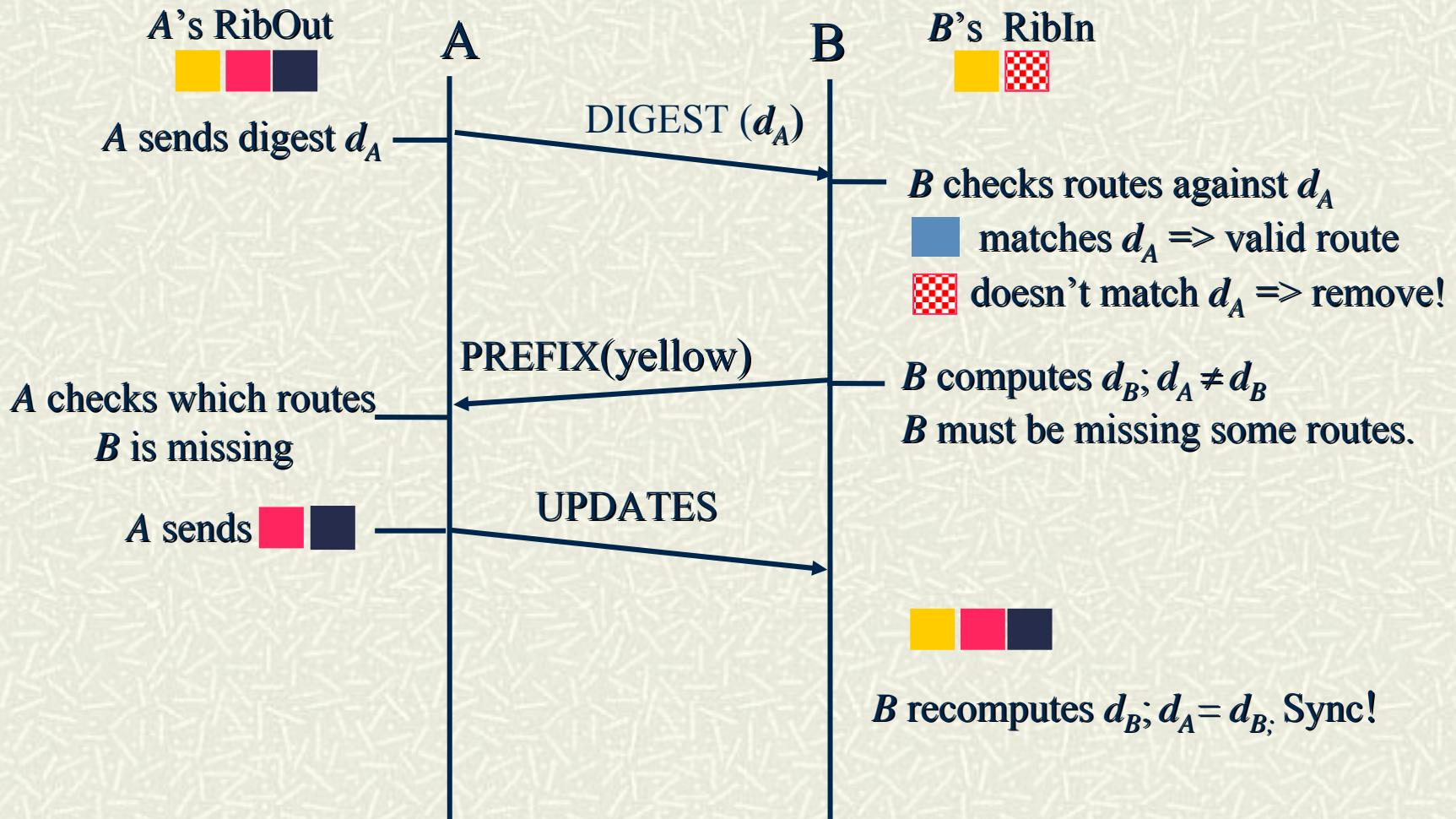
Why use Bloom Filter?

- Compact size \Rightarrow low bandwidth cost
- Low false positive rate
 - a *small* probability f that an invalid route will be considered valid.
 - $f = (1 - e^{-k \cdot n/l})^k$ (e.g. $l/n = 8$, $k = 3 \Rightarrow f = 3.1\%$)

How to remove the false positives?

- Good news: some false positives may be exposed *as the state set evolves*.
- For the remaining false positives, take advantage of the periodic digests.
 - Use seeded hash functions, e.g. salted MD5;
 - Change the seed periodically;
- Residual error $e = f^i$ ($i = \#$ different seeds)
 - Example: 99.997% of the inconsistencies will be removed when $f = 3.1\%$ and $i = 3$.

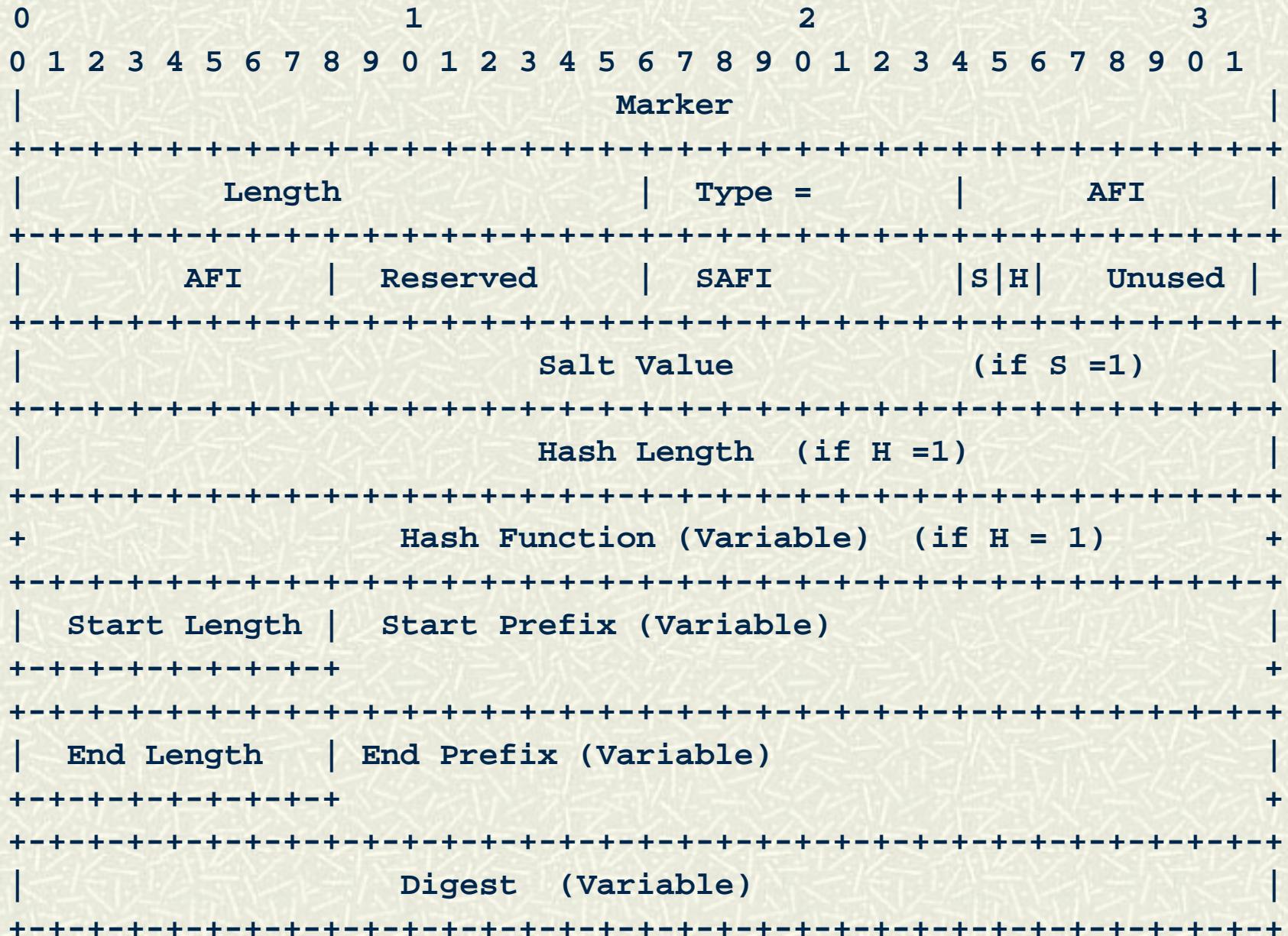
Routing Table Recovery Procedure



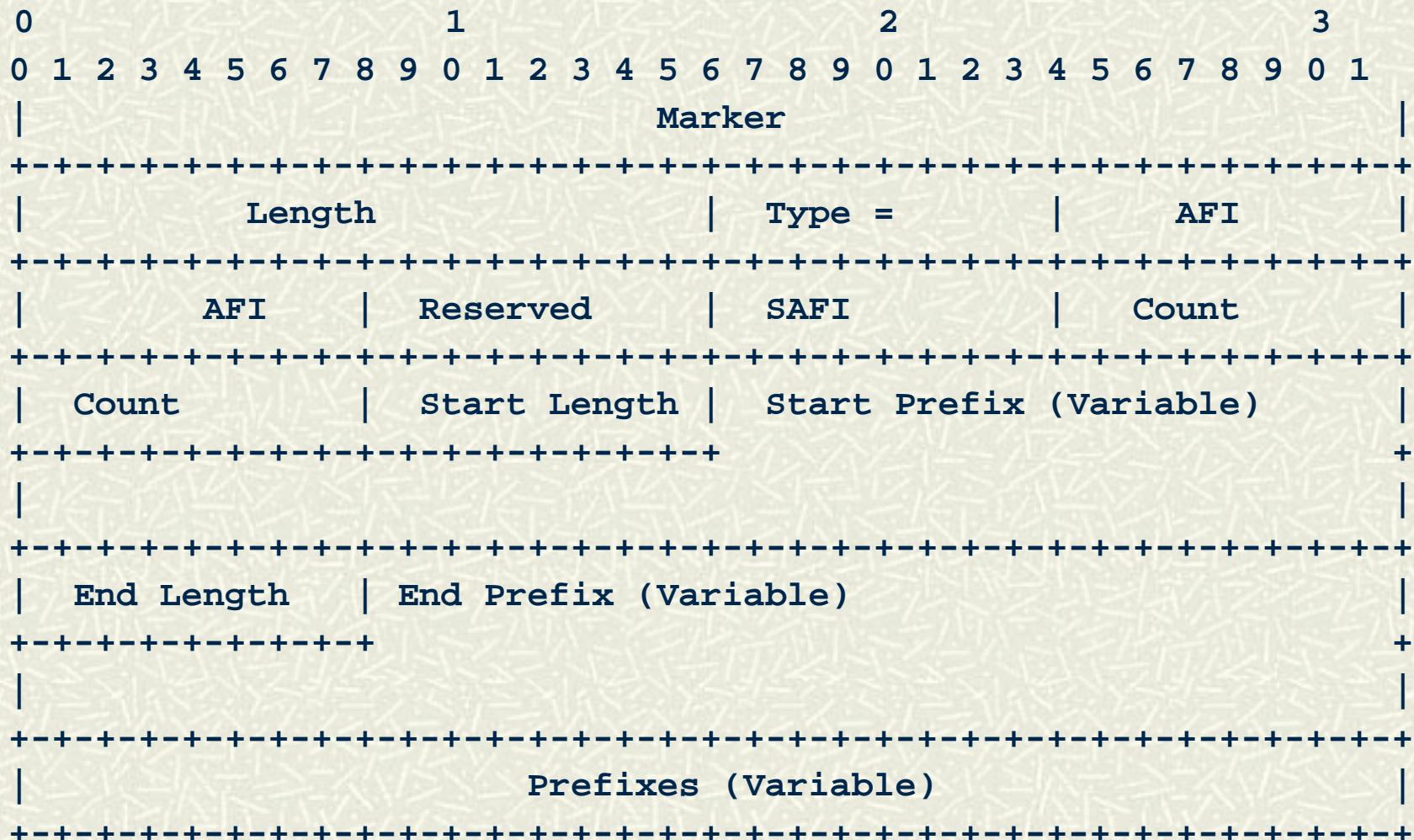
BGP OPEN Message – FRTR Negotiation

+-----+ Capability Code 'X' (1 octet) to be assigned by IANA +-----+
+-----+ Capability Length = 10 (1 octet) +-----+
+-----+ AFI (2 octets) +-----+
+-----+ Reserved (1 octet) +-----+
+-----+ SAFI (1 octet) +-----+
+-----+ FRTR interval (4 octets) +-----+
+-----+ FRTR digest size (2 octets) +-----+

Digest Message



Prefix Message

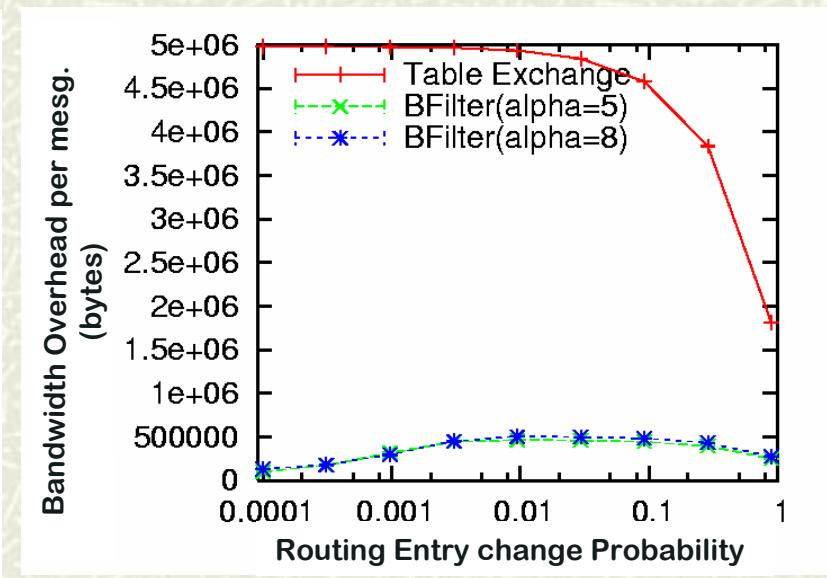
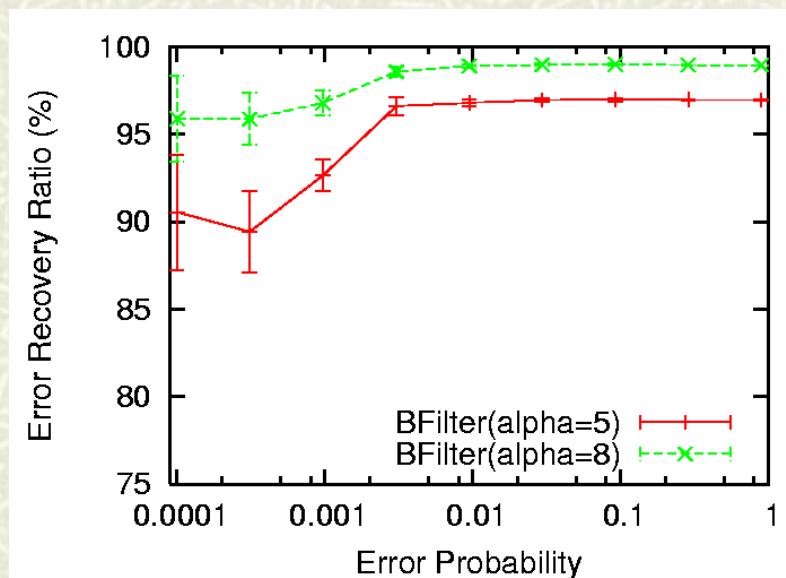


Evaluation of FRTR

- Data source: routing tables from 11 ISPs
- Methodology
 - Inject random errors into each routing table;
 - Simulate two BGP peers: one with a original table and the other with a corrupted table.
- Bloom Filter Parameters
 - 3 hash functions
 - Encoding ratio: 5 and 8

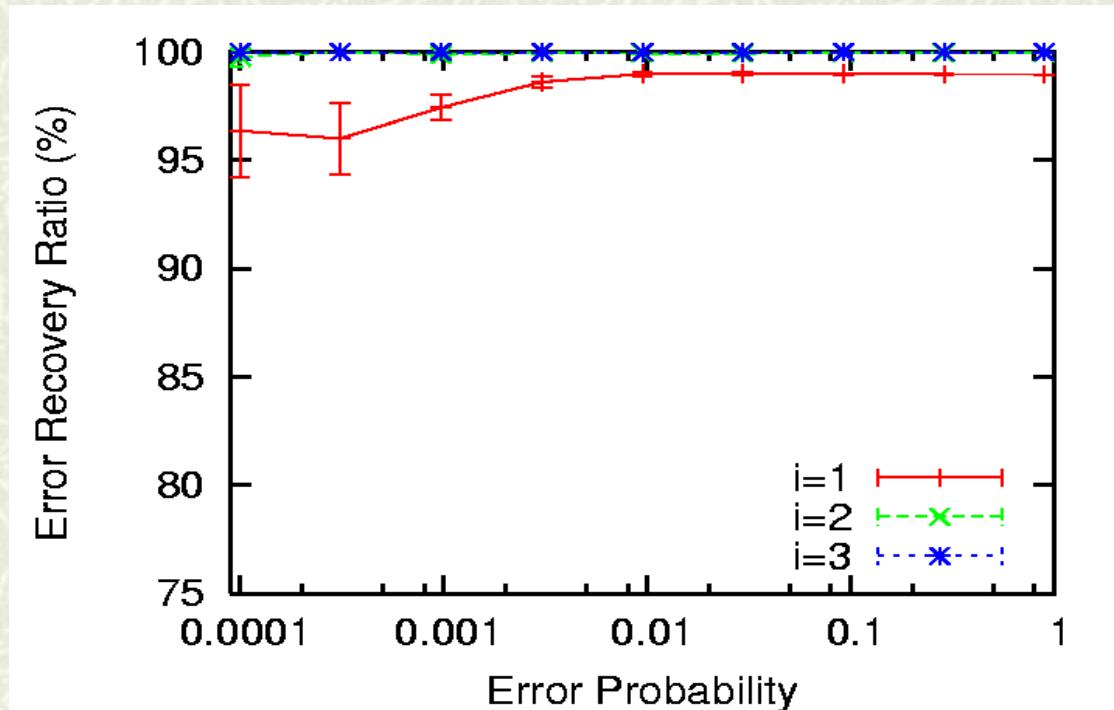
FRTR Performance

- Error recovery ratio: 91%~97% for encoding ratio of 5, 96% ~ 99% for encoding ratio of 8
- Bandwidth overhead per message (encoding ratio = 5): as low as 75K bytes, at most 460K bytes



Performance of Salted Digests

- Error recovery ratio (encoding ratio = 8): close to 100% after a few seed changes.



Summary

- FRTR encodes routing table to efficiently detect and recover unnoticeable errors.
- FRTR significantly reduces the overhead of routing table recovery after session resets.