#### Transparent TCP Timestamps draft-scheffenegger-tcpm-timestamp-

negotiation-03

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Timestamp Echo on SYN TCP RTO calculation Changes since -02 draft





RFC1323 states, a sender SHOULD set TSecr to zero, if no valid timestamp is known

 Investigated this issue against Waikato Dataset (IP Header trace) covering 1. Jan. 2005 – 17. Aug. 2005

- 1 016 245 629 [SYN] packets in trace
  - 878 264 323 (86.4%) unique [SYN]
- 3 634 [SYN] && TSecr != 0 (3.6 ppm)
  - 3 613 (99.4%) unique
  - 586 unique host pairs
  - 457 unique senders
  - 1 953 unique TSecr values

- Majority of TSecr values sent on [SYN] is seen in at least one preceeding packet – always from a different TCP session between the two hosts
- 1 783 (of 1 953; 91.3%) [SYN] && TSecr != 0 preceeded by
  - 1 191 [FIN, ACK]
  - 5 [FIN, PSH, ACK]
  - 573 [ACK]
  - 13 [PSH, ACK]
- Remaining 171 [SYN] have no packet from opposite side (unidirectional flow captured)

passive OS signature on [SYN] && TSecr != 0

	MSS SACK TS NOP WS	3520 (96.86%)
•	NOP NOP TS	56(1.54%)
•	MSS NOP WS NOP NOP TS	22(0.61%)
•	MSS NOP NOP TS	16(0.44%)
•	MSS NOP NOP TS WS	9(0.25%)
•	MSS SACK TS	8 ( 0.22%)
	MSS NOP WS NOP NOP TS NOP	1(0.03%)
•	MSS NOP WS NOP NOP TS NOP NOP	1(0.03%)
•	MSS NOP WS NOP NOP TS NOP WS	1(0.03%)

- passive OS signature on [SYN] && TSecr != 0
  - p0f fails to identify senders (stringent heuristics)
  - SinFP uses "relaxed" heuristics
    - 1 IPv4: HEURISTIC0/P2: BSD: Darwin: 8.6.0
    - 2 IPv4: BH0FH0WH1OH0MH1/P2: BSD: FreeBSD: 4.4
    - 4 IPv4: HEURISTIC0/P2: BSD: FreeBSD: 4.10
    - 12 IPv4: BH0FH0WH0OH0MH1/P2: BSD: Darwin: 8.6.0

106 IPv4: unknown

4 IPv4: BH0FH0WH1OH0MH1/P2: GNU/Linux: Linux: 2.2.x 5 IPv4: BH0FH0WH0OH1MH0/P2: GNU/Linux: Linux: 2.4.x 5 IPv4: BH0FH0WH1OH0MH1/P2: GNU/Linux: Linux: 2.4.x 6 IPv4: BH0FH0WH1OH0MH0/P2: GNU/Linux: Linux: 2.4.x 49 IPv4: BH0FH0WH0OH0MH1/P2: GNU/Linux: Linux: 2.6.x

409 IPv4: BH0FH0WH0OH0MH1/P2: GNU/Linux: Linux: 2.4.x 2814 IPv4: HEURISTIC0/P2: GNU/Linux: Linux: 2.4.x

- 3509 sessions initiated from Linux (96.5%)
- 106 sessions initiated from unknown Host OS
- 19 sessions identified as BSD derived
- Known feature of Linux, when either sysctl\_tcp\_tw\_reuse or sysctl\_tcp\_tw\_recycle are enabled. This is used to shorten the TimeWait interval before sockets may be reused between two hosts.
- Observed TSecr values are not random, but preceeded in all observable instances by identical TSval (in a different TCP session), clearly indicating some kind of per-host timestamp cache.

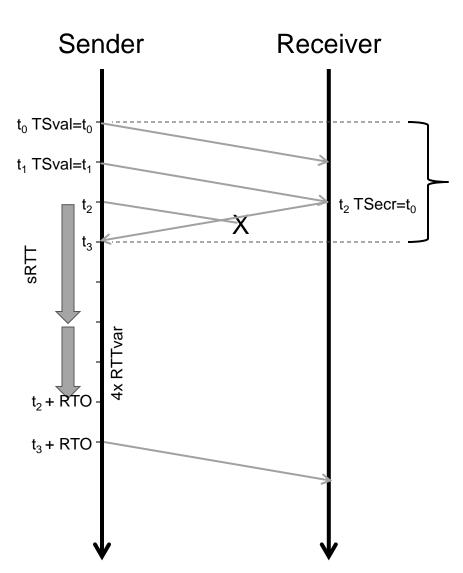
- As timestamp values represented uptime in some way, values are not random distributed. Would lead to reduced false positive negotiation with current proposal.
- Most TSecr values are seen in [FIN] segments (61.3%). As [FIN] packets are closing a TCP session, small semantic change to send TSval = 0 on [FIN] would further alleviate this issue. As with timestamp offset randomization (draft-gont-tcpm-tcp-timestamps), may break linux feature - revert to regular timewait session closure.

## TCP RTO calculation

Defined in RFC6298: RTO = sRTT + max(G, k \* RTTvar) ; k = 4 RTTvar =  $(1-\beta)$  \* RTTvar +  $\beta$  \* | sRTT - R' | ;  $\beta$  = 1/4 sRTT =  $(1-\alpha)$  \* sRTT +  $\alpha$  \* R' ;  $\alpha$  = 1/8

- The sampled RTT (R') no longer includes delACK processing variability
- RTTvar will become smaller
- RTT represents primarily represents network delay
- One RTT extra delay before RTO expires added by restarting the RTO timer for each seen ACK

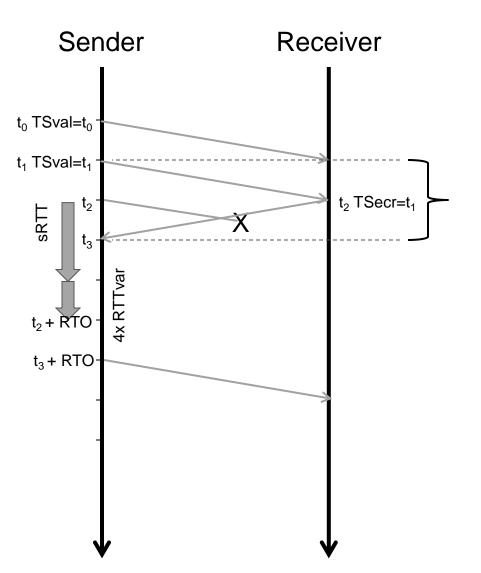
## TCP RTO calculation – RFC1323 / 6298



- RTT includes delayed ACK delay
- Retransmission permitted after t<sub>2</sub> + RTO

Typically at t<sub>3</sub> + RTO

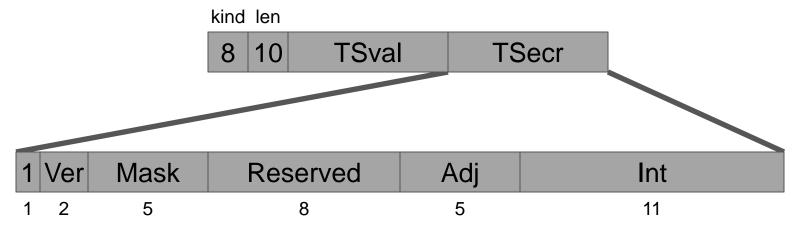
## **TCP RTO calculation – new semantic**



- RTT excludes delayed ACK delay
- Retransmission permitted after
  t<sub>2</sub> + RTO – too fast?

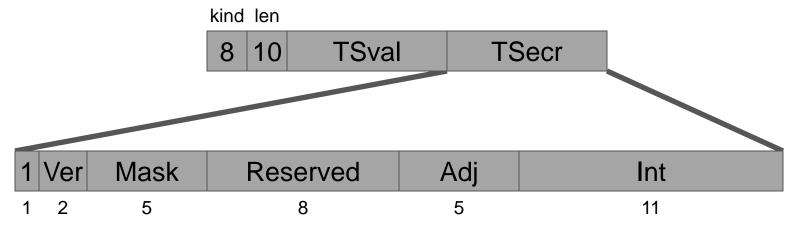
Explicitly stipulate max(t<sub>2</sub>,t<sub>3</sub>) + RTO (last received ACK + RTO)?

# Major Change since -01



- Redefined lower 16 bits
- No longer "alike" IEEE 754 floating point representation
  - No "special case" handling when evaluating Interval
  - Interval now signaled as scale/value pair like TCP window scaling
  - Conceptually, the Interval is a large integer, right-shifted "Adj" times to fit into "Int" with the most significant bit
  - Derived calculations stay identical (i.e. OWD) as with previous draft
  - Allows for implicit "clock source" quality signaling by leading zero bits in "Int"

# Major Change since -01

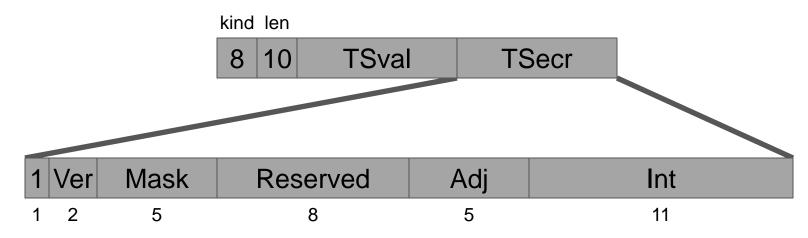


- Definition: 1 sec = 0x40 0000 0000 (39 bit length)
  - Right shifted 28 times, to fit into 11 bit
  - Int = 0x400, Adj=28

Alternative representations:

- Int=0x200, Adj=29;
- Int=0x100, Adj=30;
- Int=0x080, Adj=31
- Allows for clock tick intervals between ~16 sec and a few nanoseconds, each with up to 10 alternative representations (clock source quality levels)

# Major Change since -01



 Expected to be (compile time) static value for current slow running TCP clock sources (100 µs or longer intervals)

# Transparent TCP Timestamps

Questions:

- semantic change to timestamp to have TSval set to zero on [FIN]?
- Changed signaling of Interval (simple shift vs. IEEE-754 like float) less problematic?
- Ready for adoption as a WG Item?





#### Thank you!

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