



# MPLS Fast Re-route using extensions to LDP

draft-kini-mpls-frr-ldp-02

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IETF 82, Taipei, NOV 2011

# Solution

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- › Recovery characteristics similar to RSVP-TE FRR (sub 50msec recovery) for traffic on routed paths
- › Local repair mechanism – Only PLR reacts to the failure trigger to recover the traffic
- › Backup shortest path (BSP) LDP LSP setup before failure whenever LFA does not exist
- › 100% coverage for link, node and SRLG failure
- › BSP LSP starts at PLR and merges into shortest path LDP LSP tree. Merge point referred to as BSP-MP

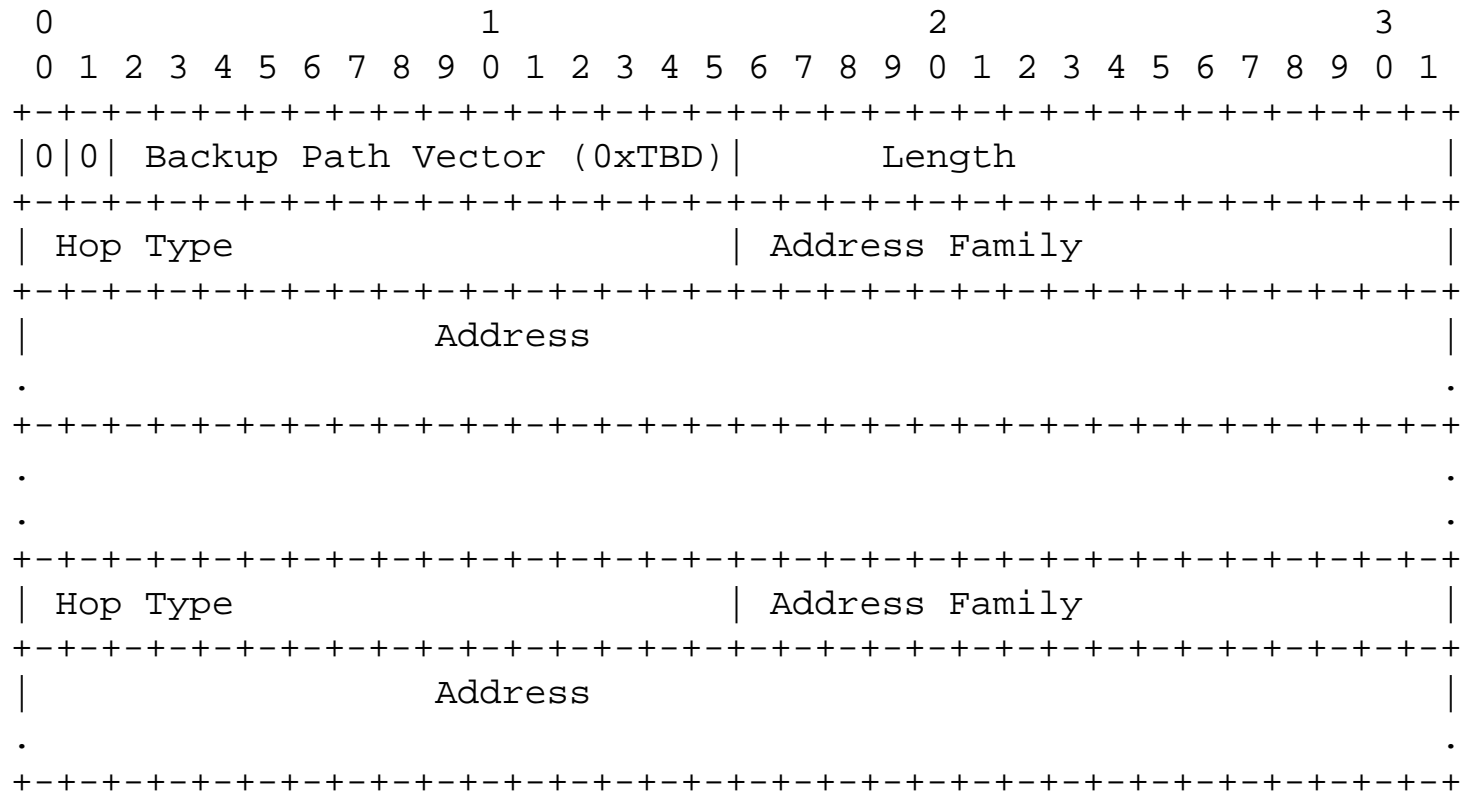
# Diff 01 and 02

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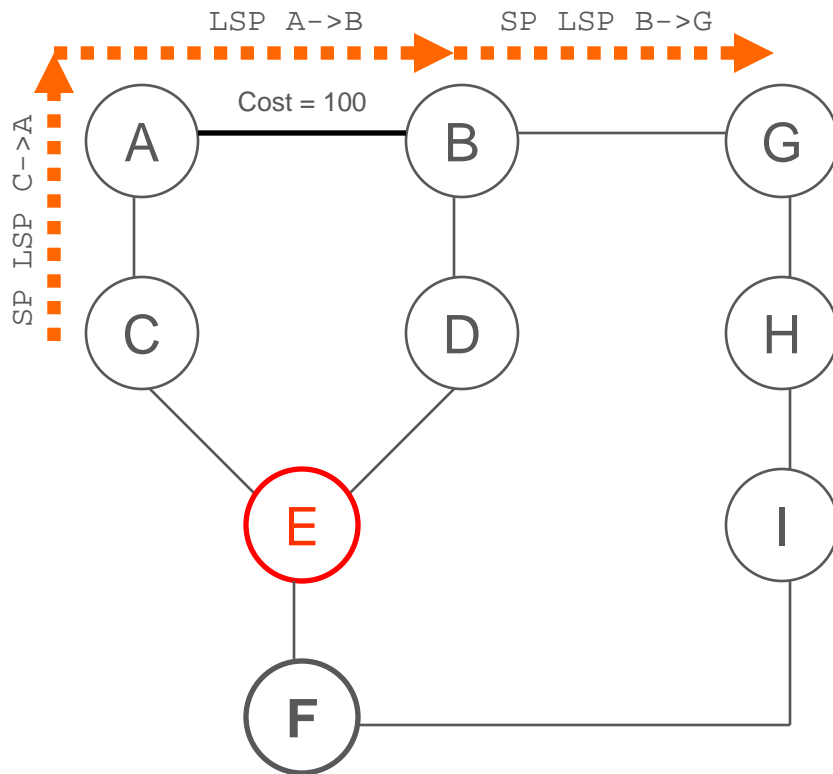
- › Simplified BSP LSP definition
  - BSP LSP is Shortest path LSP from PLR to the BSP Merge point
  - When there is no single shortest path, multiple shortest path LSPs are stitched together
  
- › Backup Path Vector TLV differentiates BSP LSP related label exchanges from other FEC label exchanges
  
- › Failure Element TLV and Tunneled FEC TLV are removed
  
- › Defined capability TLVs

# Backup Path Vector TLV

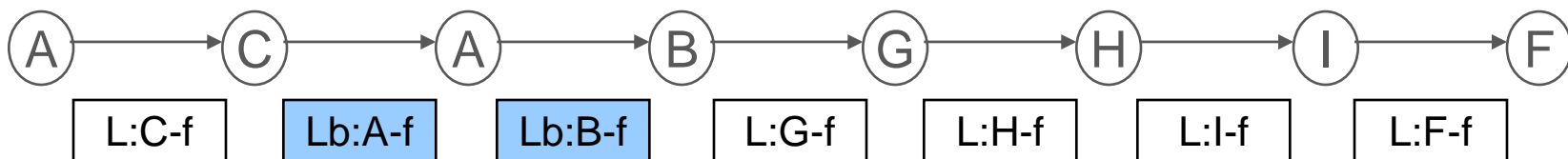
- › Backup Path Vector TLV – contains list of addresses of stitching LSRs



# Example



- › Cost of link A-B is 100, rest is 1
- › Failure of node E
- › C is PLR
- › BSP LSP = C – A – B – G
- › L:X-y – Label assigned by Node X for FEC y
- › Lb:X-y – BSP Label assigned by Node X for FEC y



# Operational details

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- › Per-nexthop protection can reduce number of BSP LSPs
  
- › What happens when a single shortest-path LSP is not available for tunneling ?
  - Shortest path LSPs are stitched together using simple extensions to LDP
  
- › Protocol Extensions
  - Backup Path Vector TLV

# Comparison with other approaches

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## › LDP over RSVP

- Less OpEx (managing one less protocol). Simplicity.
- Less protocol state
- Multi-path on backup

## › LFA & Not-via

- Full coverage
- Re-uses MPLS FRR infrastructure
- Simple

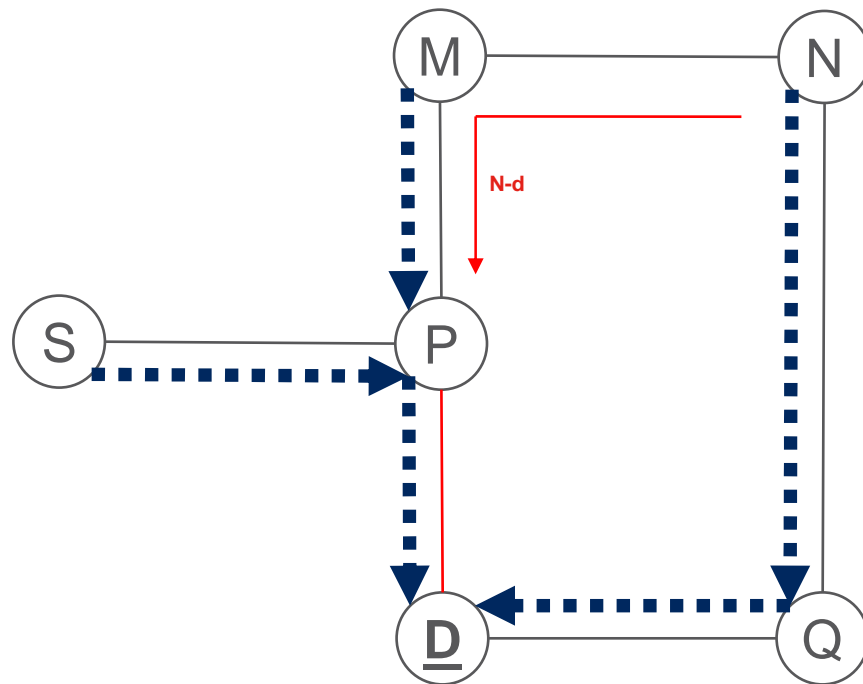
# Questions/Comments

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(Backup slides included with several examples)



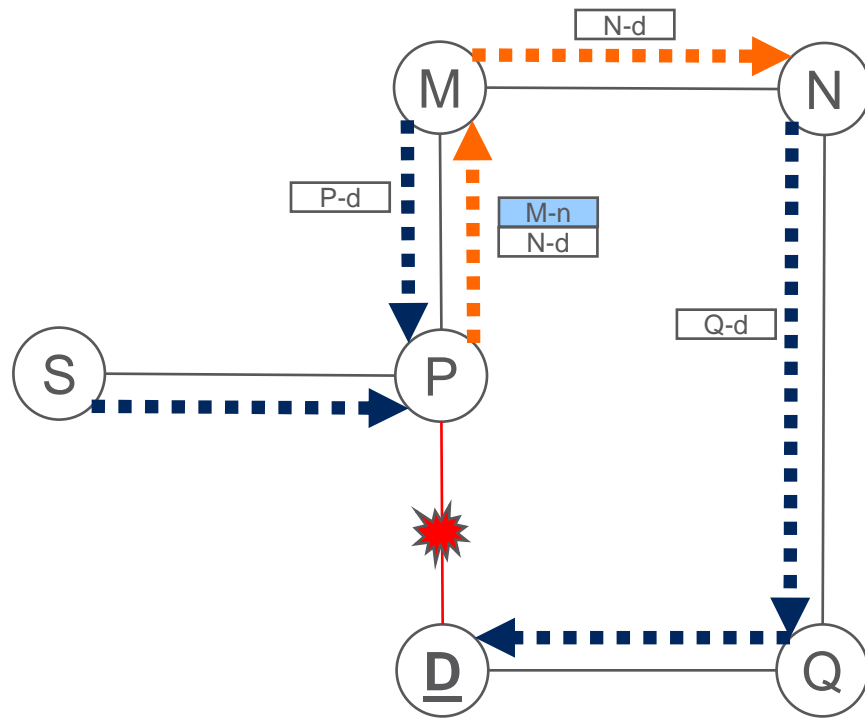
# Link failure protection example



- › Protect link P-D failure
- › For Destination D
  - P is PLR
  - N is merge point
  - N advertises label **N-d** to P for the backup shortest-path LSP
  - **N-d** is the shortest-path LDP LSP label at N for D
  - P uses shortest-path LSP from P to N to tunnel label **N-d**

.....➔ Traffic flow over shortest path LSP

# Link failure protection fast re-routed traffic



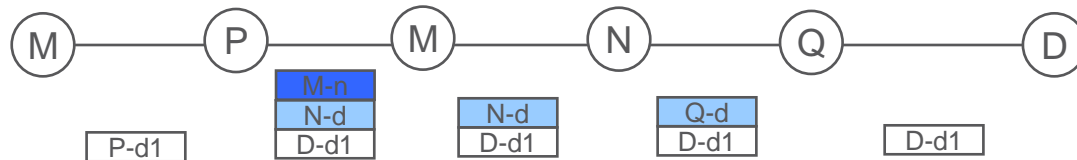
Fast re-routed traffic

FRR traffic paths to D when link P-D fails

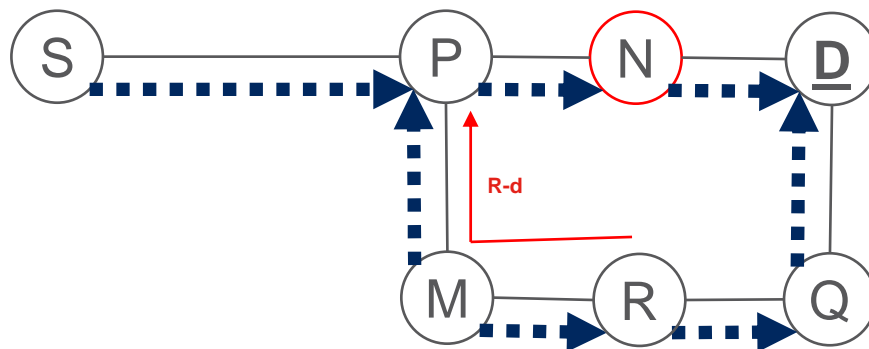
- › P, M, N, Q, D
- › S, P, M, N, Q, D
- › M, P, M, N, Q, D

For entire network

- › No 'new' labels needed in the network
- › 12 additional label advertisements needed



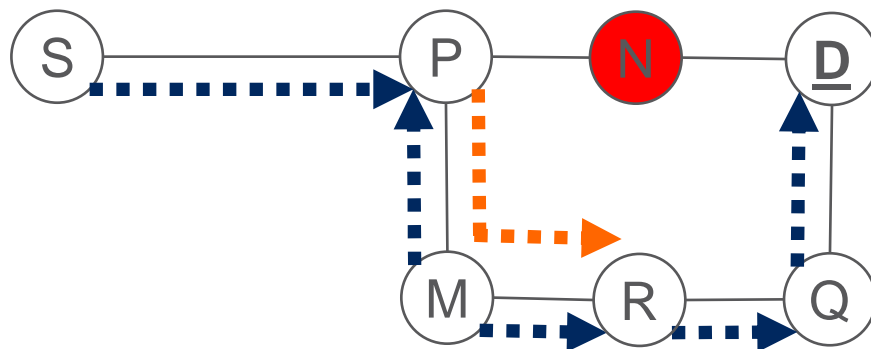
# Node failure protection example



- › Node N failure
- › Destination D
- › P is PLR
- › R is merge point
- › R advertises label **R-d** to P for the backup shortest-path LSP

.....➔ Traffic flow over shortest path LSP

# Node failure protection fast re-routed traffic



FRR traffic paths to D  
when node N fails

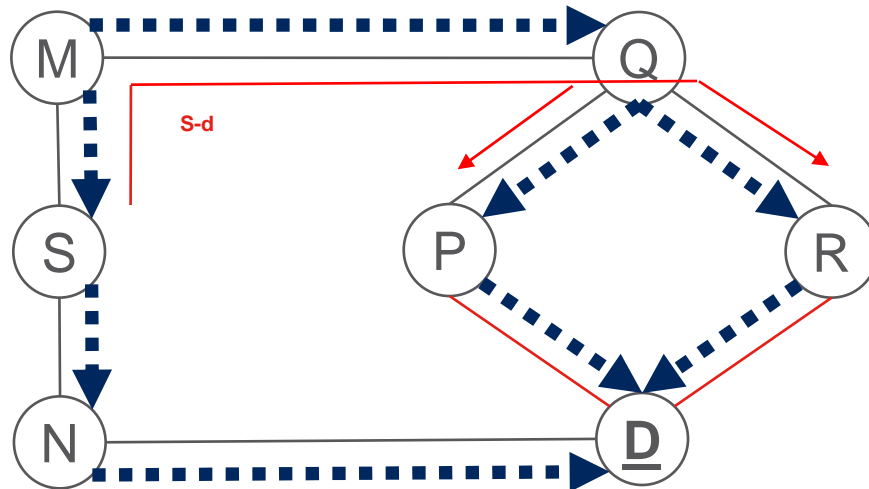
- › P, M, R, Q, D
- › S, P, M, R, Q, D
- › M, P, M, R, Q, D

For entire network

- › No 'new' labels needed in the network
- › 6 additional label advertisements needed

.....➔ Fast re-routed traffic

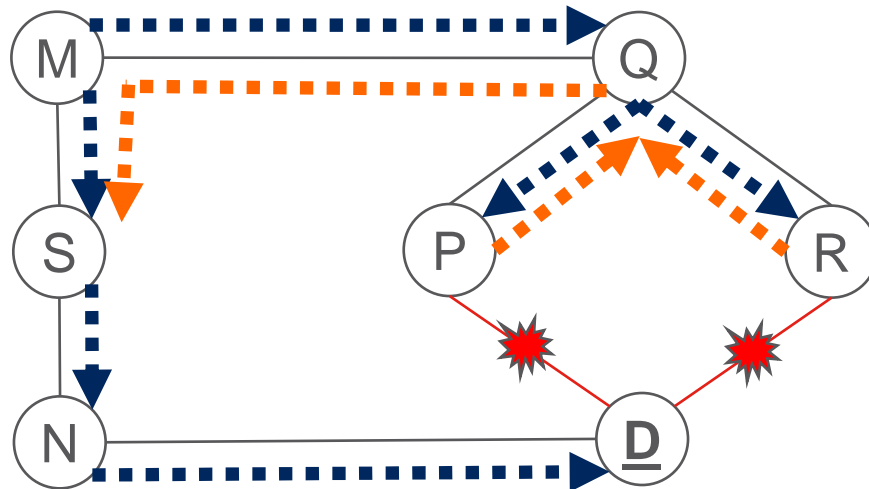
# SRLG failure protection example



- › SRLG (link P-D, link R-D) failure
- › Destination D
- › P, R are PLRs
- › S is merge point
- › S advertises its shortest path LSP label (S-d) to P and R for failure against SRLG

.....➔ Traffic flow over shortest path LSP

# SRLG failure protection fast re-routed traffic

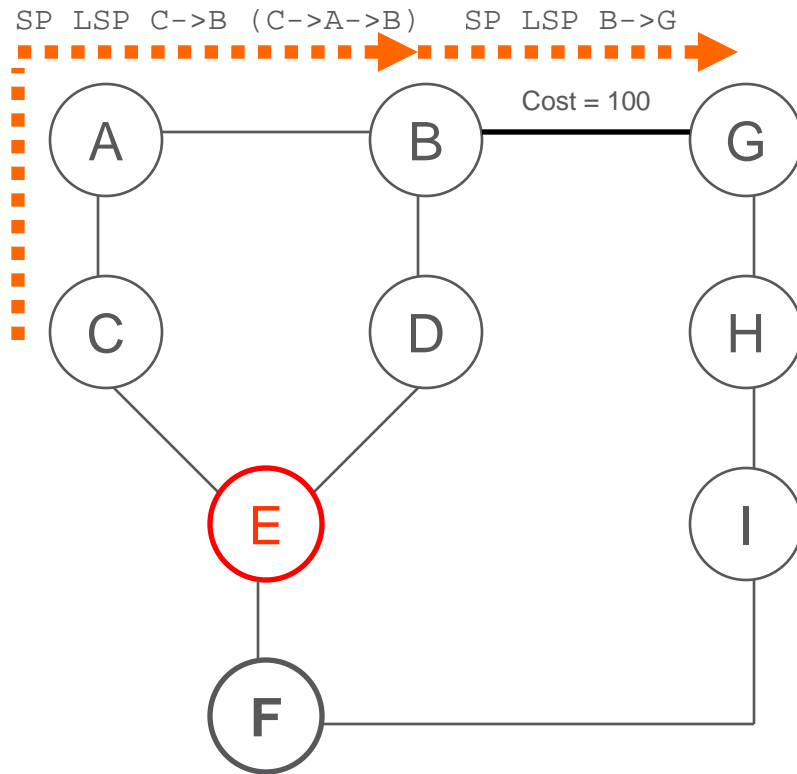


FRR traffic paths to D  
when SRLG fails

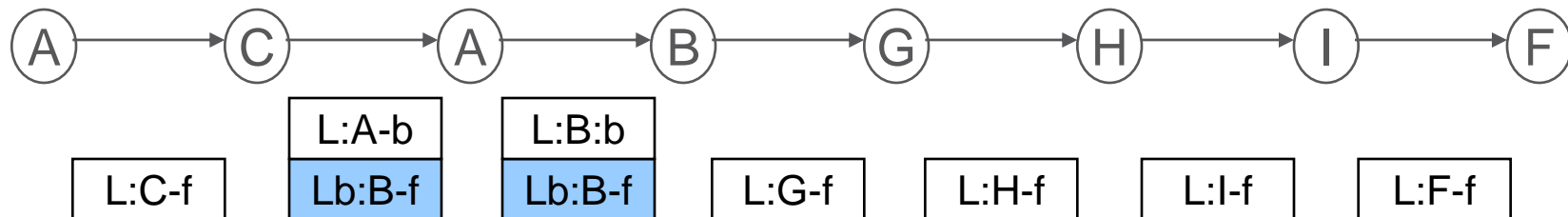
- > P, Q, M, S, N, D
- > Q, P, Q, M, S, N, D
- > Q, R, Q, M, S, N, D
- > M, Q, P, Q, M, S, N, D
- > M, Q, R, Q, M, S, N, D

.....> Fast re-routed traffic

# Example - 2



- › Cost of link B-G is 100, rest is 1
- › Failure of node E
- › C is PLR
- › BSP LSP = C – B – G
- › L:X-y – Label assigned by Node X for FEC y
- › Lb:X-y – BSP Label assigned by Node X for FEC y



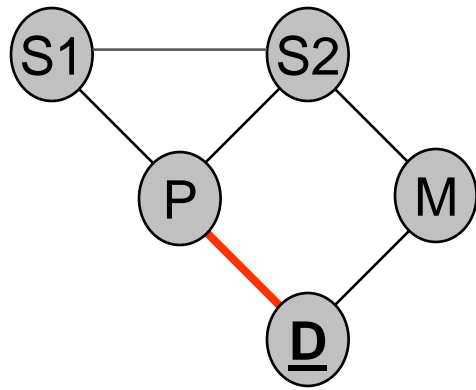
# Computation

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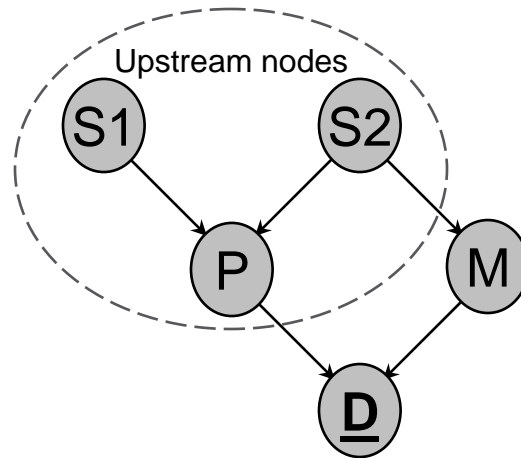
- › SPT for a destination
- › Failure at PLR
- › Nodes upstream of failure in the SPT is affected
- › Nodes not upstream of failure in the SPT is not affected
- › Compute SPT with “failure” excluded – Exclude-SPT
- › Alternate path from PLR to destination in Exclude-SPT merges back into SPT @ BSP-MP (not upstream of failure)
- › BSP LSP from PLR to BSP-MP protects the traffic under failure



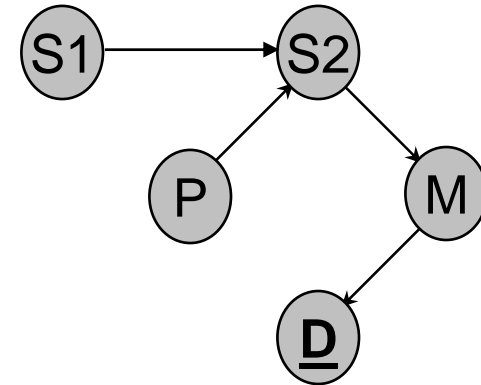
# SPT & Exclude-SPT



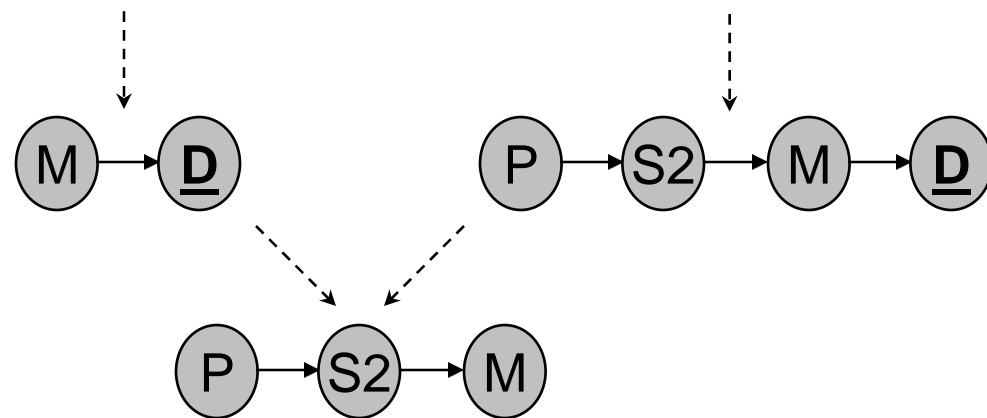
Topology



SPT



Exclude-SPT



BSP LSP