MPLS Fast Re-route using extensions to LDP

draft-kini-mpls-frr-ldp-01

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Motivation and Goal

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> LDP LSPs are widely deployed.
(IGPastrest 50th) ec recovery for traffic on routed paths

(IGP shortest path)

> Full coverage needed

Solution should be self-contained. It should be independent of other protocols and mechanisms such as IP-FRR,

Solution characteristics

- Computation intensive tasks are performed much before the actual failure (during steady state).
- Only PLR reacts to the failure trigger to recover the traffic
- Actions at the PLR to recover the traffic are simple (and precomputed)

Solution summary

>

Defined for link-state IGP. And for platform label space.
Backup shortest path (BSP) LDP LSP setup before failure

whenever LFA does not exist

BSP LSP starts at PLR and merges into shortest path LDP

LSP tree. Merge point referred to as BSP-MP.

- Fast re-route action on detecting failure
- -PLR label switches to pre-selected BSP LDP LSP
 - Stack label to aggregate failures. Use shortest-path LSP from PLR

Computation

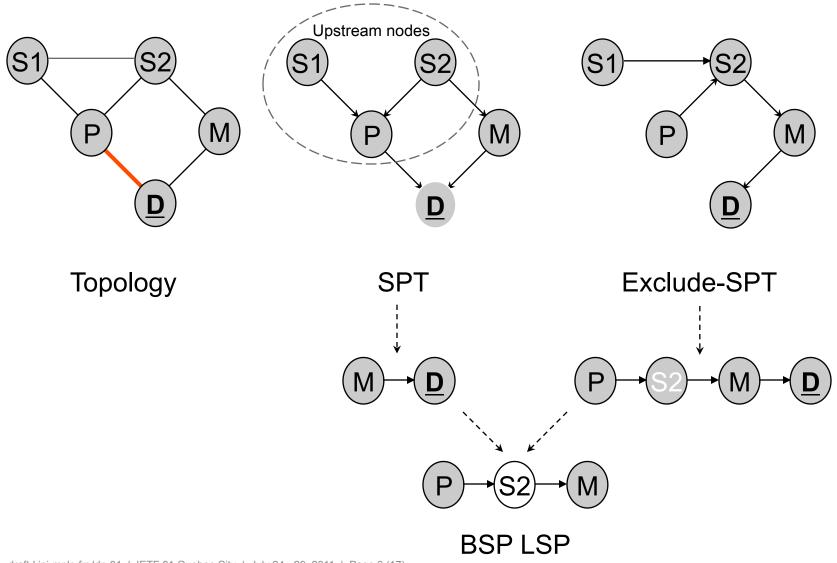
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- SPT for a destination
- > Failure at PLR
- Nodes upstream of failure in the SPT is affected
- Nodes not upstream of failure in the SPT is not affected
- Compute SPT with "failure" excluded Exclude-SPT

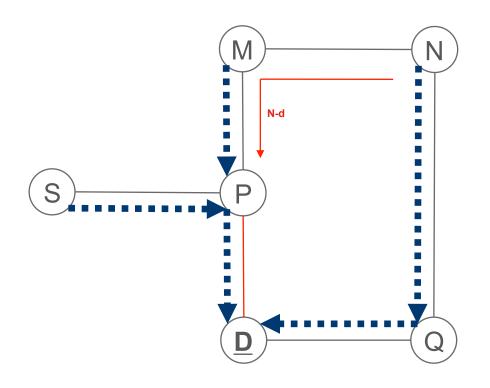
Alternate path from PLR to destination in Exclude-SPT merges back into SPT @ BSP-MP (not upstream of failure)

BSP LSP from PLR to BSP-MP protects the traffic under

SPT & Exclude-SPT



Link failure protection example

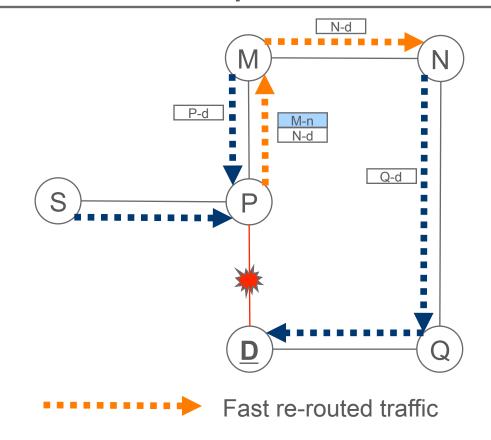


Protect link P-D failureFor Destination **D**

- N advertises label for the backup shortest-path for the backup shortest-path LSP
- N-d is the shortest-path LDP LSP label at N for D
- P uses shortest-path LSP from P to N to tunnel labelN-d

Traffic flow over shortest path LSP

Link failure protection fast re-routed traffic

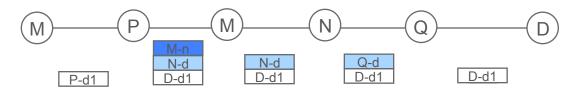


FRR traffic paths to **D** when link P-D fails

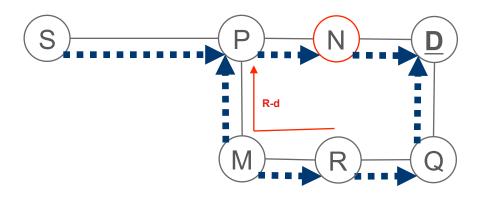
- > P, M, N, Q, **D**
- > S, P, M, N, Q, <u>D</u>
- > <u>M, P, M, N, Q, D</u>

For entire network

- No 'new' labels needed in the network
- 12 additional label advertisements needed



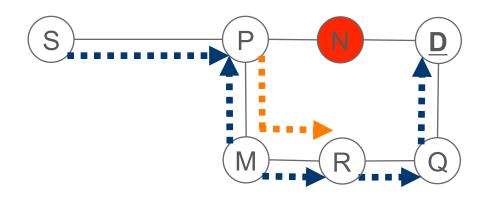
Node failure protection example



- Node N failure
- Destination <u>D</u>
- > P is PLR
- > R is merge point
- R advertises label R-d to P for the backup shortest-path LSP

Traffic flow over shortest path LSP

Node failure protection fast re-routed traffic



FRR traffic paths to **D** when node N fails

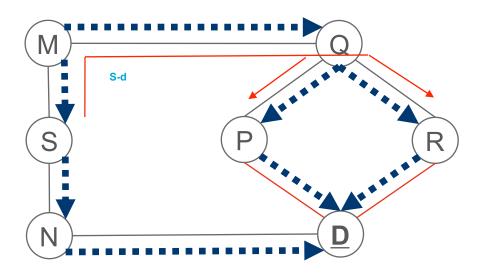
- > P, M, R, Q, **D**
- > S, P, M, R, Q, **D**
- » M, P, M, R, Q, **D**

For entire network

- No 'new' labels needed in the network
- 6 additional label advertisements needed



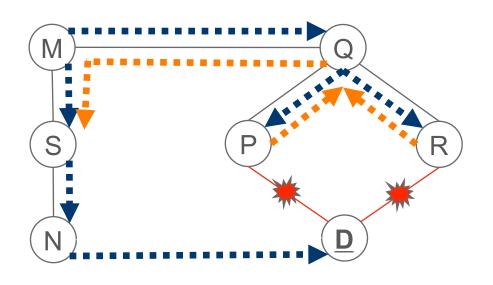
SRLG failure protection example



- S advertises its shortest SBFG (link **PD**
- > Destination __ apainst Republic
- S is merge point
- S advertises its shortest path LSP label (S-d) to P and R for failure against SRLG

Traffic flow over shortest path LSP

SRLG failure protection fast re-routed traffic



FRR traffic paths to **D** when SRLG fails

- > P, Q, M, S, N, **D**
-) Q, P, Q, M, S, N, **D**
- > Q, R, Q, M, S, N, **D**
- > M, Q, P, Q, M, S, N, **D**
- » M, Q, R, Q, M, S, N, **D**

Fast re-routed traffic

Operational details

- > Per-nexthop protection can reduce number of BSP LSPs
- What happens when a shortest-path LSP is not available for tunneling?
 - Explicit routing for BSP LSP using extensions to LDP
- > Protocol Extensions
 - Failure Element TLV
 - -Tunneled FEC TLV (when label stacking not used)
 - -Backup Path Vector TLV

Comparison with other approaches

> LDP over RSVP

- Less OpEx (managing one less protocol). Simplicity.
- Less protocol state
- Multi-path on backup

> LFA & Not-via

- Full coverage
- -Re-uses MPLS FRR infrastructure
- No IP address management issues

Future Work

Analyze applicability

Questions/Comments