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Distributed Mobility Anchoring
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Abstract

This document defines the mobility management solutions in the context of a distributed mobility management deployment. It consider the problem of assigning a mobility anchor at the initiation of a flow. In addition, the mid-session switching of the mobility anchor in a distributed mobility management environment is considered.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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This Internet-Draft will expire on April 17, 2016.

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1. Introduction

A key requirement in distributed mobility management [RFC7333] is to enable traffic to avoid traversing single mobility anchor far from the optimal route. Recall that distributed mobility management solutions do not make use of centrally deployed mobility anchor [Paper-Distributed.Mobility]. As such, a flow SHOULD be able to have its traffic changing from traversing one mobility anchor to traversing another mobility anchor as the mobile node moves, or when changing operation and management (OAM) requirements call for mobility anchor switching, thus avoiding non-optimal routes. This draft proposes distributed mobility anchoring solutions.

The needs of IP-layer mobility support are diverse so that the use of distributed anchoring may differ according to the needs.

A mobile node (MN) may be running a flow with its correspondent node (CN) for which the source IP address of this flow belongs to MN's network. That is, it is anchored to an access router (anchor) belonging to MN's network. When there are multiple anchors, the flow may need to select the anchor when it is initiated (Section 3). Using an anchor in MN's network has the advantage that the packets can simply be forwarded according to the forwarding table. The anchor may be in the MN's network when the flow was initiated. As

the MN moves from one network to another, IP address no longer belongs to the new network. To order that the IP address of the flow is in the new network different methods can be used dependent on the needs of the flow. If the ongoing IP flow can cope with an IP prefix/address change, the flow can be reiniated with a new IP address anchored in the new network (Section 3.1). On the other hand, if the ongoing IP flow cannot cope with such change, the IP address anchoring can be moved from the original network to the new network (Section 3.2).

2. Conventions and Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL","SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

All general mobility-related terms and their acronyms used in this document are to be interpreted as defined in the Mobile IPv6 base specification [RFC6275], the Proxy Mobile IPv6 specification [RFC5213], and the DMM current practices and gap analysis [RFC7429]. This includes terms such as mobile node (MN), correspondent node (CN), home agent (HA), home address (HoA), care-of-address (CoA), local mobility anchor (LMA), and mobile access gateway (MAG).

In addition, this document uses the following term:

Home network of an application session (or of an HoA): the network that has allocated the IP address (HoA) used for the session identifier by the application running in an MN. An MN may be running multiple application sessions, and each of these sessions can have a different home network.

IP prefix/address anchoring: An IP prefix, i.e., Home Network Prefix (HNP), or address, i.e., Home Address (HoA), allocated to a mobile node is topologically anchored to a node when the anchor node is able to advertise a connected route into the routing infrastructure for the allocated IP prefix.

Internetwork Location Management (LM) function: managing and keeping track of the internetwork location of an MN. The location information may be a binding of the IP advertised address/prefix, e.g., HoA or HNP, to the IP routing address of the MN or of a node that can forward packets destined to the MN. It is a control plane function.

In a client-server protocol model, location query and update messages may be exchanged between a Location Management client (LMc) and a Location Management server (LMs).

With separation of control plane and data plane, the LM function is in the control plane. It may be a logical function at the control plane node, control plane anchor, or mobility controller.

It may be distributed or centralized.

Forwarding Management (FM) function: packet interception and forwarding to/from the IP address/prefix assigned to the MN, based on the internetwork location information, either to the destination or to some other network element that knows how to forward the packets to their destination.

This function may be used to achieve indirection. With separation of control plane and data plane, FM may split into a FM function in the data plane (FM-DP) and a FM function in the control plane (FM-CP).

FM-DP may be distributed with distributed mobility management. It may be a function in a data plane anchor or data plane node.

FM-CP may be distributed or centralized. It may be a function in a control plane node, control plane anchor or mobility controller.

Security Management (SM) function: The security management function controls security mechanisms/protocols providing access control, integrity, authentication, authorization, confidentiality, etc. for the control plane and data plane.

This function resides in all nodes such as control plane anchor, data plane anchor, mobile node, and correspondent node.

3. IP prefix/address anchored in current network of attachment

The IP prefix/address at the MN’s side of a flow may be anchored at the access router to which the MN is attached.

For example, when an MN attaches to a network (Net1) or moves to a new network (Net2), it is allocated an IP prefix from that network. It configures from this prefix an IP address which is typically a dynamic IP address. It then uses this IP address when it a flow is initiated. Packets to the MN in this flow are simply forwarded according to the forwarding table.

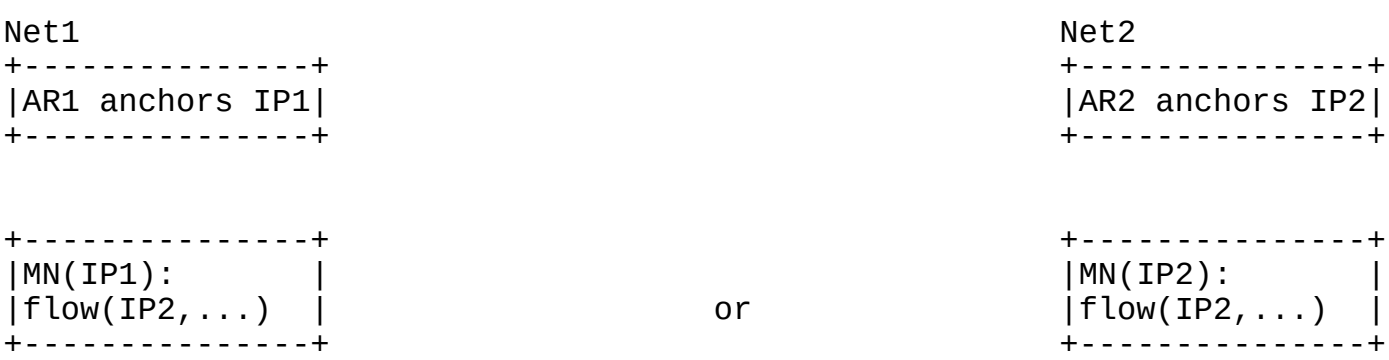

Figure 1. IP prefix/address anchored in network of attachment. MN is attached to AR1 in Net1 where it has initiated a flow using IP1 or has moved to AR2 in Net2 where it initiates a new flow using IP2.

There may be multiple IP prefixes/addresses to choose from. They may be from the same access network or different access networks. The network may advertise these prefixes with cost options [I-D.mccann-dmm-prefixcost] so that the mobile node may choose the one with the least cost. In addition, these IP prefixes/addresses may be of different types regarding whether mobility support is needed [I-D.dmm-ondemand-mobility-api]. A flow will need to choose the appropriate one according to whether it needs IP mobility support.

With on-demand mobility, IP mobility support is provided only when needed instead of being provided by default.

3.1. Changing to the new IP prefix/address

A straightforward choice of mobility anchoring is for a flow to use the IP prefix of the network to which the MN is attached when the flow is initiated [I-D.seite-dmm-dma]. This is shown in Figure 2.

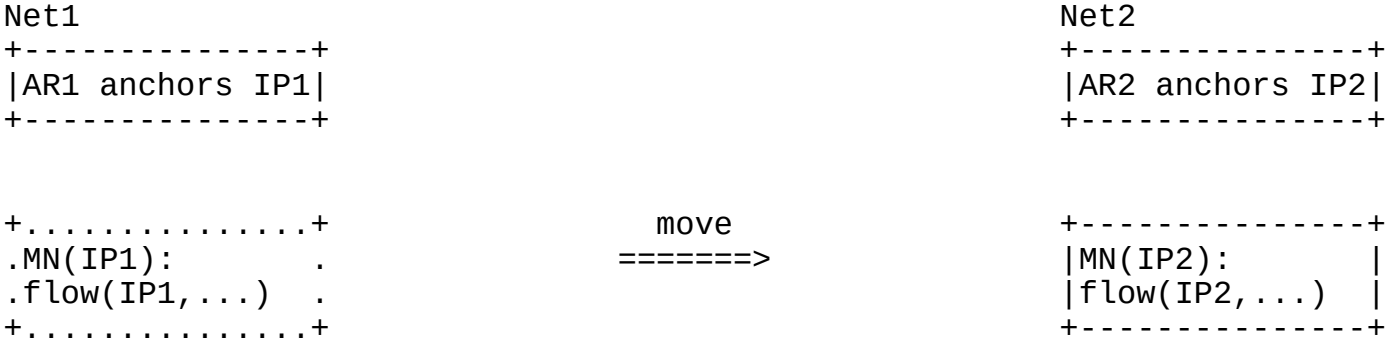

Figure 2. Changing to the new IP prefix/address. MN running a flow using IP1 in Net1 changes to running a flow using IP2 in Net2.

When IP mobility is not provided to a specific flow, the flow may use a new IP address acquired from a new network as the MN moves to the new network.

Regardless of whether IP mobility is needed, if the flow has terminated before the MN moves to a new network, the flow may subsequently restart using the new IP address allocated from the new network.

When session continuity is needed, even if a flow is ongoing as the MN moves, it may still be desirable for the flow to change to using the new IP prefix configured in the new network. The flow may then close and then restart using a new IP address configured in the new network. Yet such a change in flow may be using a higher layer mobility support which is not in the scope of this document to change the IP address of the flow.

In Figure 2, a flow initiated while the MN was in Net1 has terminated before the MN moves to a new network Net2. After moving to Net2, the MN uses the new IP prefix anchored in Net2 to start a new flow. The packets may then be forwarded without requiring IP layer mobility support.

The call flow is outlined in Figure 3.

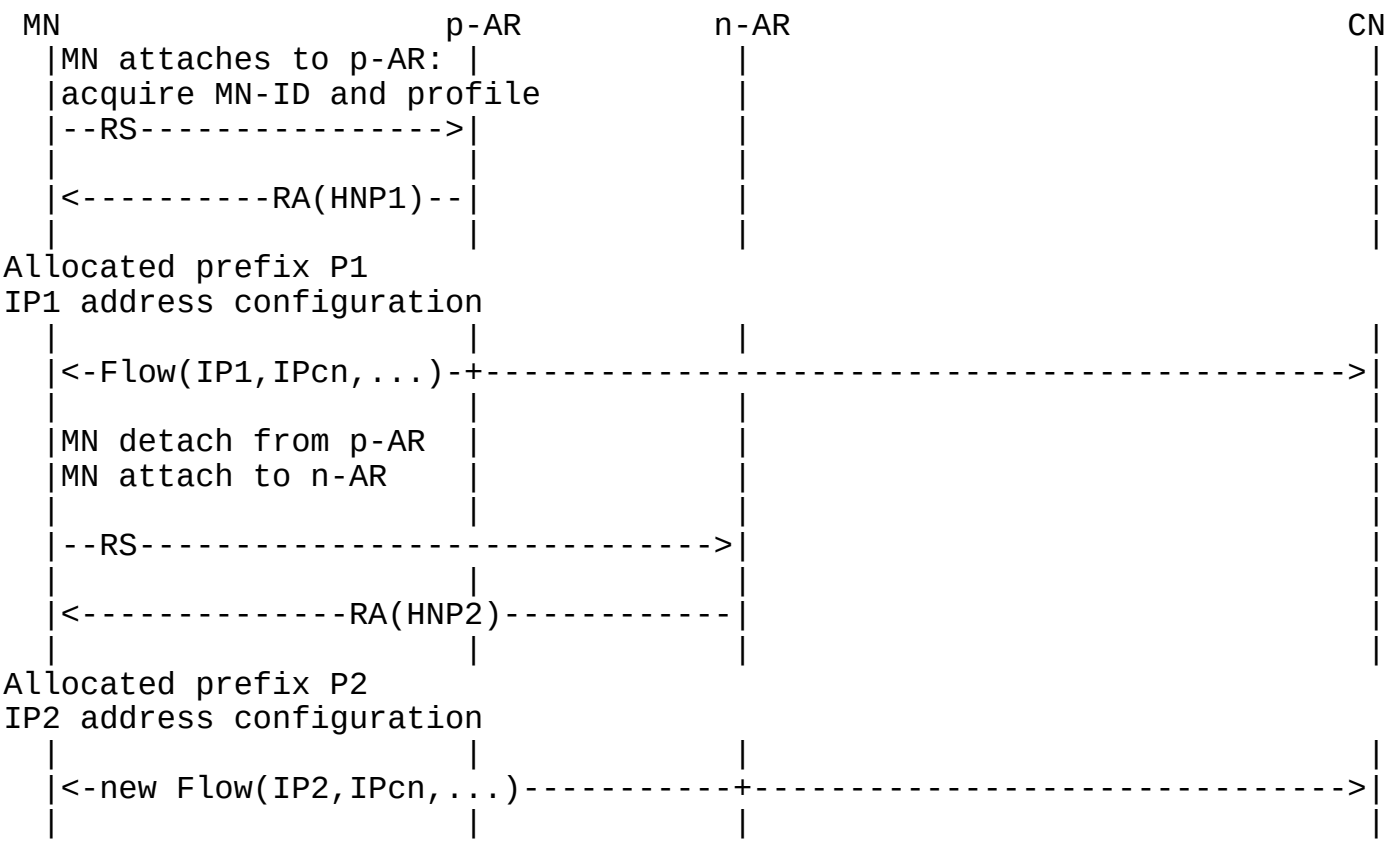

Figure 3. A flow uses the IP allocated from the network at which the MN is attached when the flow is initiated.

The security management function in the anchor node at a new network must allow to assign a valid IP prefix/address to a mobile node.

When IP mobility is needed for a flow, the mobility support may be provided by moving the IP address anchoring to the new network to be described in Section 3.2 or by using other mobility management methods ([Paper-Distributed.Mobility.PMIP] and [Paper-Distributed.Mobility.Review]) Then the flow may continue to use the IP prefix from the prior network. Yet some time later, the flow of a certain user application may be closed. If the spplication is started again, the new flow may not need to use the prior network address to avoid having to invoke IP mobility support. This is the case when the use of a permanent IP prefix/address is not needed. The flow may then use the new IP prefix in the network where the flow is initiated. Routing is again kept simpler without employing IP mobility and will remain so as long as the MN has not moved away from that network.

The call flow in this case is outlined in Figure 4.

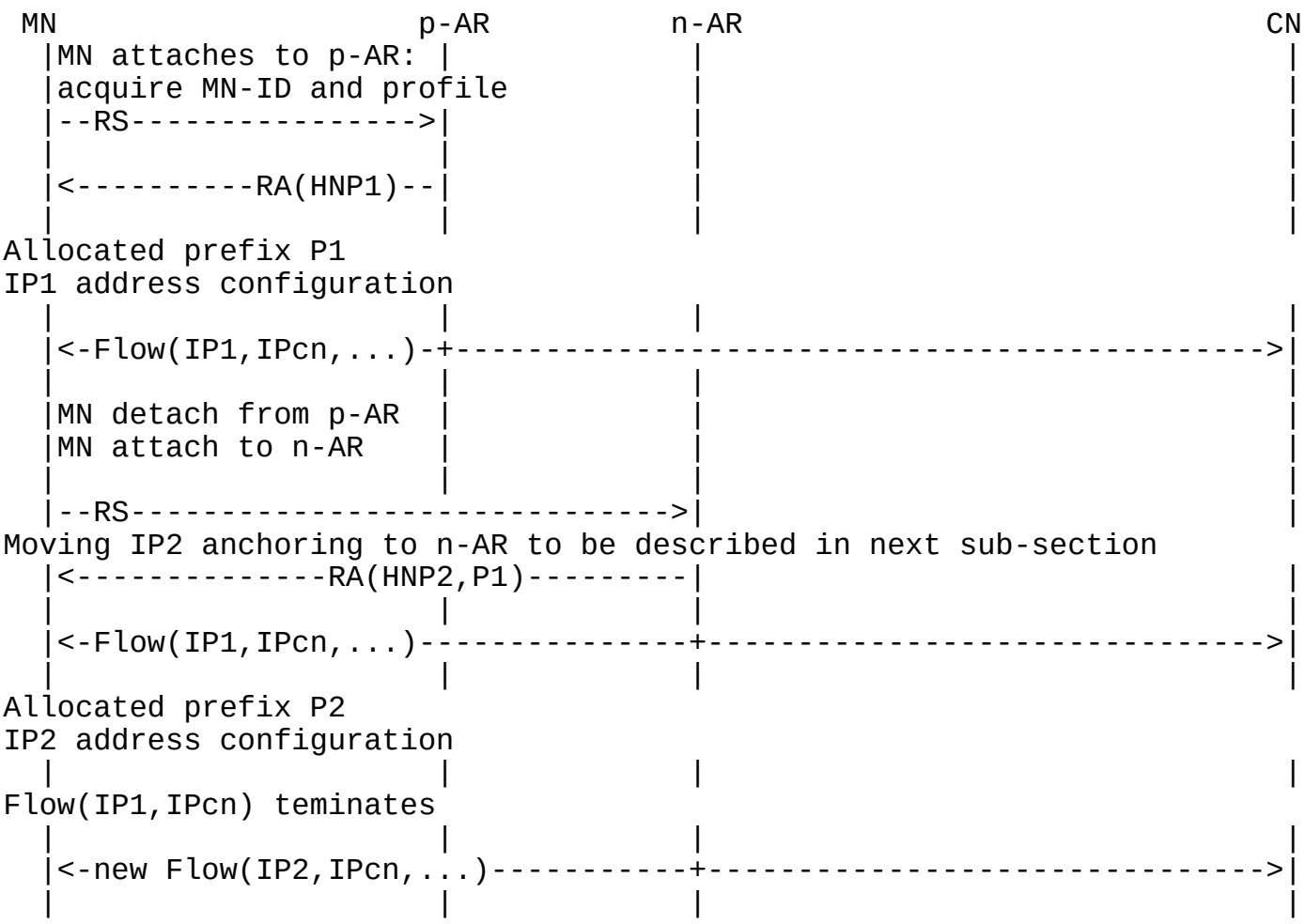

Figure 4. A flow uses the IP allocated from the network at which the MN is attached when the flow is initiated.

3.2. Moving the IP prefix/address anchor to the new network

The IP prefix/address anchor may move without changing the IP prefix/address of the flow.

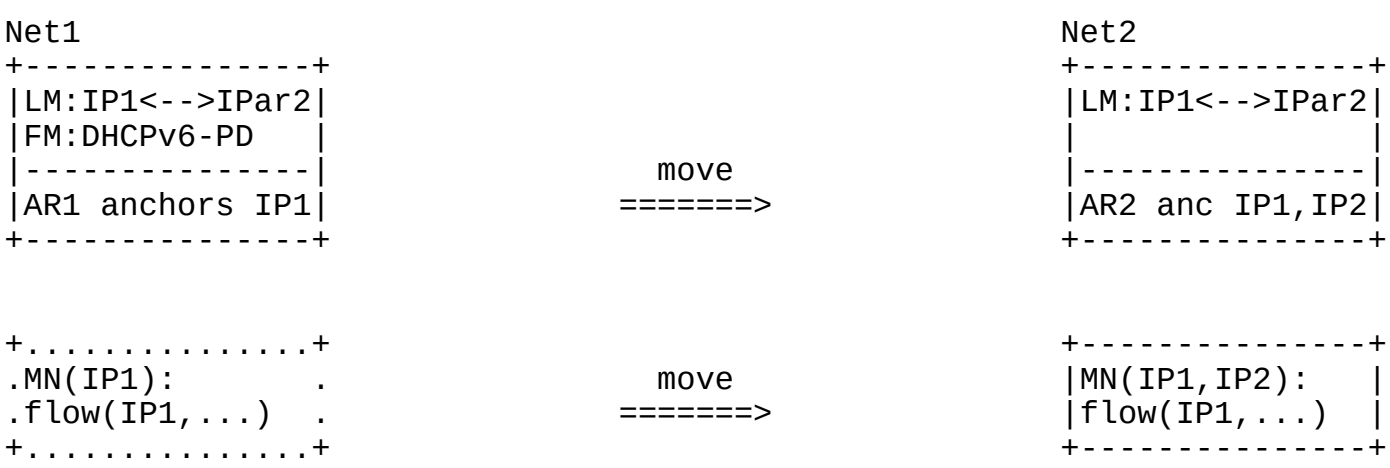

Figure 5. Moving the IP prefix/address anchor to the new network. MN with flow using IP1 in Net1 continues to run the flow using IP1 as it moves to Net2.

As an MN with an ongoing session moves to a new network, the flow may preserve session continuity by moving the original IP prefix/address of the flow to the new network. An example is in the use of BGP UPDATE messages to change the forwarding table entries as described in [I-D.mccann-dmm-flatarch] and also for 3GPP Evolved Packet Core (EPC) network in [I-D.matsushima-stateless-uplane-vepc].

The security management function in the anchor node at a new network must allow to assign the original IP prefix/address used by the mobile node at the previous (original) network. As the assigned original IP prefix/address is to be used in the new network, the security management function in the anchor node must allow to advertise the prefix of the original IP address and also allow the mobile node to send and receive data packets with the original IP address.

The security management function in the mobile node must allow to configure the original IP prefix/address used at the previous (original) network when the original IP prefix/address is assigned by the anchor node in the new network. The security management function in the mobile node also allows to use the original IP address for the previous flow in the new network.

3.2.1. Centralized control plane

An example of moving the IP prefix is in the case where Net1 and Net2 both belong to the same operator network with separation of control and data planes ([I-D.liu-dmm-deployment-scenario] and [I-D.matsushima-stateless-uplane-vepc]), where the controller may send to the switches/routers the updated information of the

forwarding tables with the IP addressing anchoring of the original IP prefix/address at AR1 moved to AR2 in the new network. That is, the IP address anchoring in the original network which was advertising the prefix will need to move to the new network. As the anchoring in the new network advertises the prefix of the original IP address in the new network, the forwarding tables will be updated so that packets of the flow will be forwarded according to the updated forwarding tables. Figure 6 shows such a case where the functions LM, FM-CP are centralized whereas the FM-DP's are distributed.

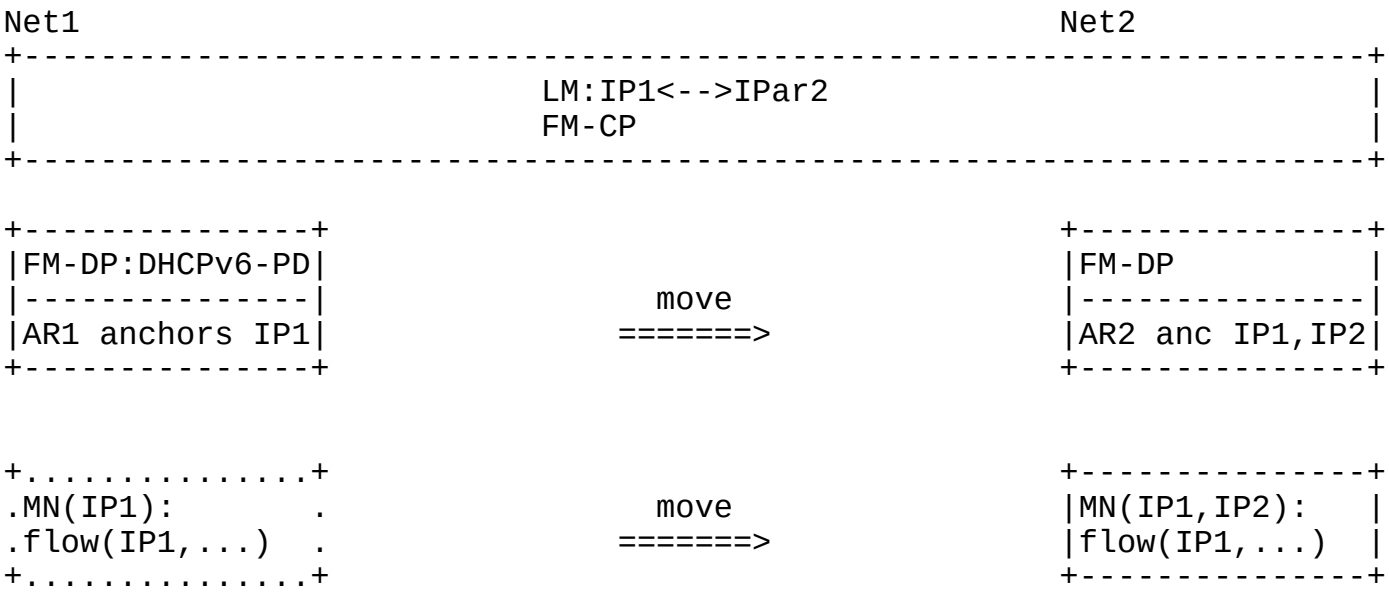

Figure 6. Moving the IP prefix/address anchor to the new network and with LM and FM-CP in a centralized control plane whereas the FM-DP's are distributed.

The call flow in Figure 7 shows that MN is allocated HNP1 when it attaches to the p-AR. A flow running in MN may or may not need IP mobility. If it does, it may continue to use the previous IP prefix. If it does not, it may use a new IP prefix allocated from the new network.

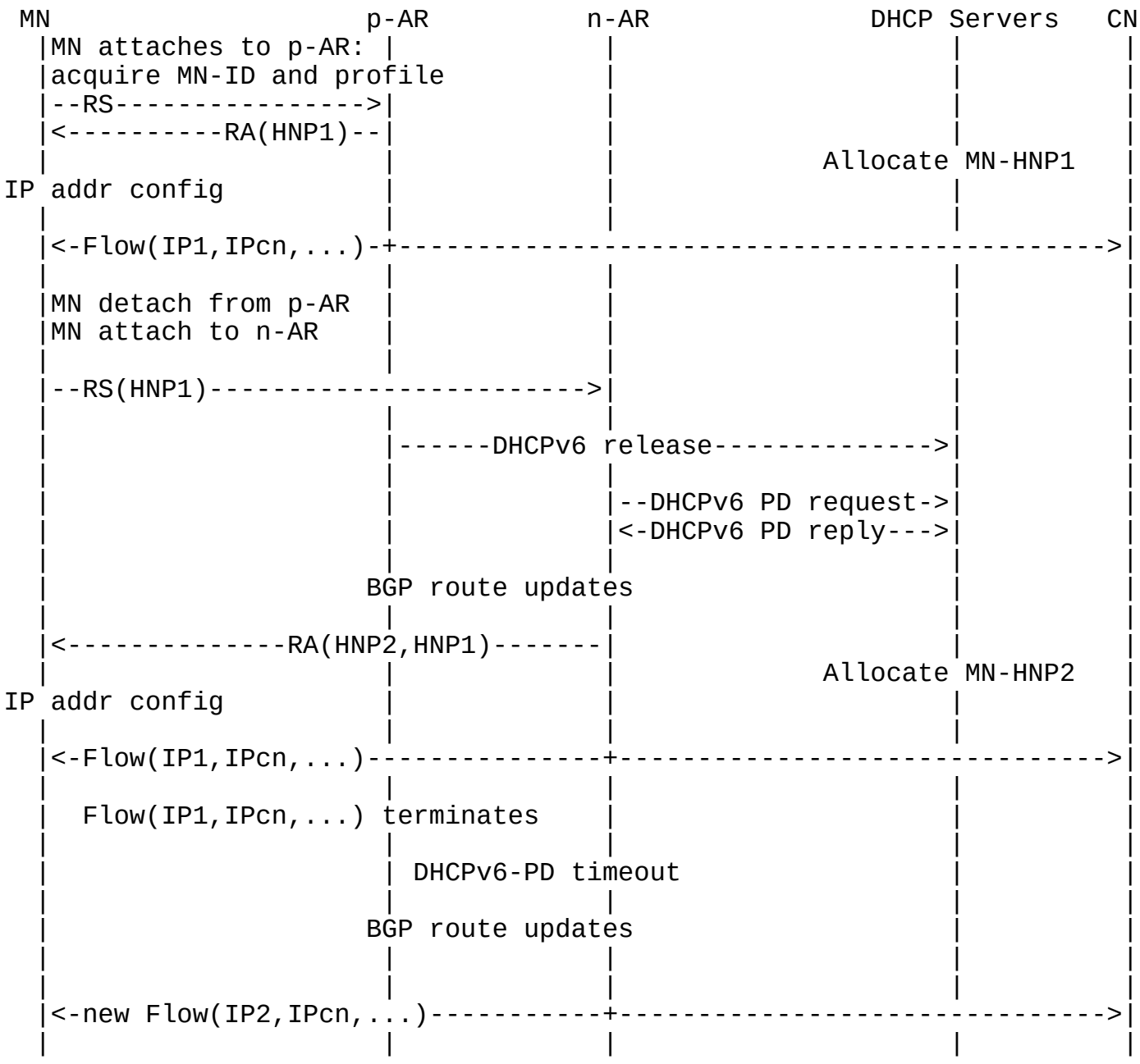

Figure 7. DMM solution. MN with flow using IP1 in Net1 continues to run the flow using IP1 as it moves to Net2.

As the MN moves from p-AR to n-AR, the p-AR as a DHCP client may send a DHCP release message to release the HNP1. It is now necessary for n-AR to learn the IP prefix of the MN from the previous network so that it will be possible for Net2 to allocate both the previous network prefix and the new network prefix. MN may provide its

previous network prefix information by including it to the RS message [I-D.jhlee-dmm-dnpp].

Knowing that MN is using HNP1, the n-AR sends to a DHCP server a DHCPv6-PD request to move the HNP1 to n-AR. The server sends to n-AR a DHCPv6-PD reply to move the HNP1. Then BGP route updates will take place here.

In addition, the MN also needs a new HNP in the new network. The n-AR may now send RA to n-AR, with prefix information that includes HNP1 and HNP2. The MN may then continue to use IP1. In addition, the MN is allocated the prefix HNP2 with which it may configure its IP addresses. Now for flows using IP1, packets destined to IP1 will be forwarded to the MN via n-AR.

As such flows have terminated and DHCP-PD has timed out, HNP1 goes back to Net1. MN will then be left with HNP2 only, which it will use when it now starts a new flow.

3.2.2. Hierarchical network

A hierarchy may also exist as shown the Figure 8. Here the IP prefix allocated to the MN is anchored at an edge router (ER) supporting multiple access routers to which the MN may be attached. Mobility of the MN involving change of AR but not of ER may be accomplished using tunneling between the ER and the AR or using some other L2 mobility mechanism.

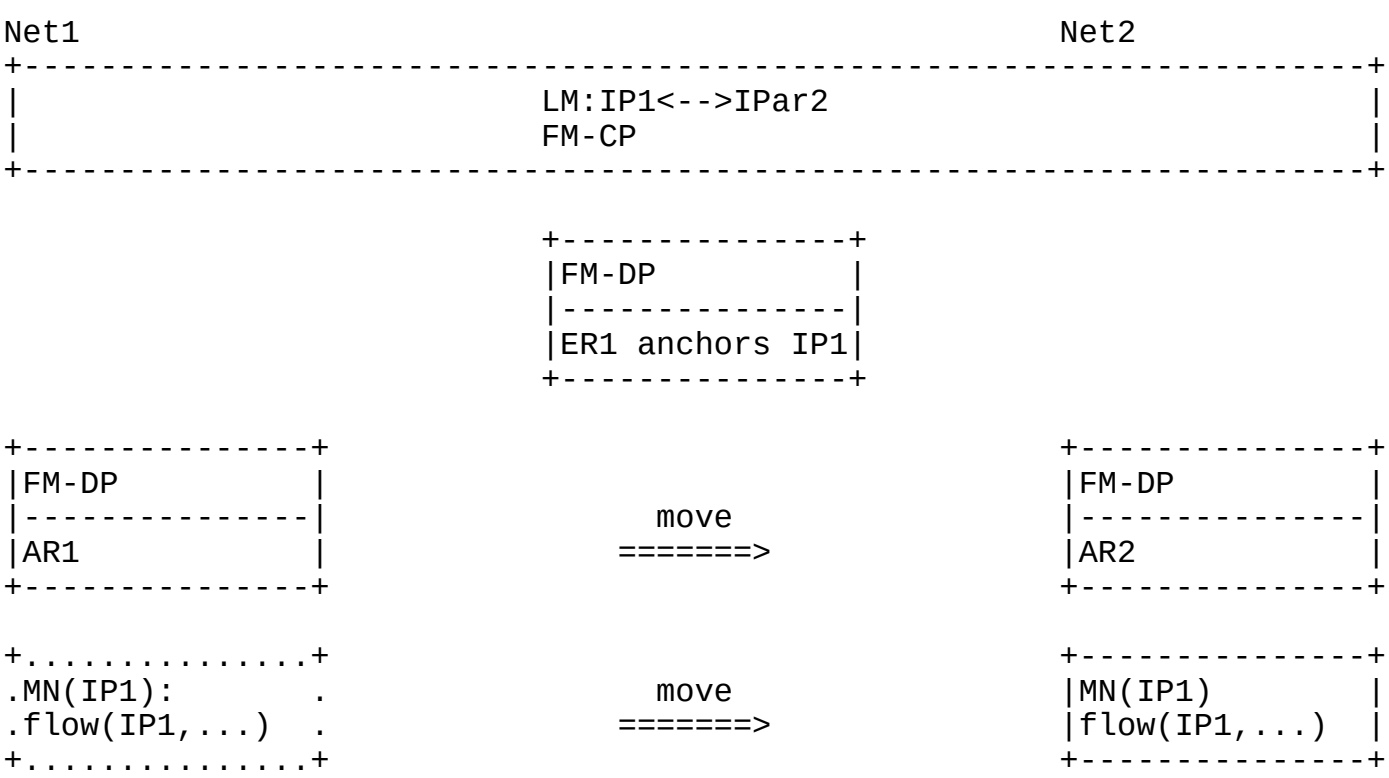

Figure 8. Mobility without involving change of IP anchoring in a network with hierarchy in which the IP prefix allocated to the MN is anchored at an Edge Router supporting multiple access routers to which the MN may connect to.

The mobility event shown in Figure 9 involves a change of the IP prefix anchoring. ER1 acting as a DHCP-PD client may exchange message with the DHCP server to release the prefix IP1. Meanwhile, ER2 acting as a DHCP-PD client may exchange message with the DHCP server to delegate the prefix IP1 to ER2.

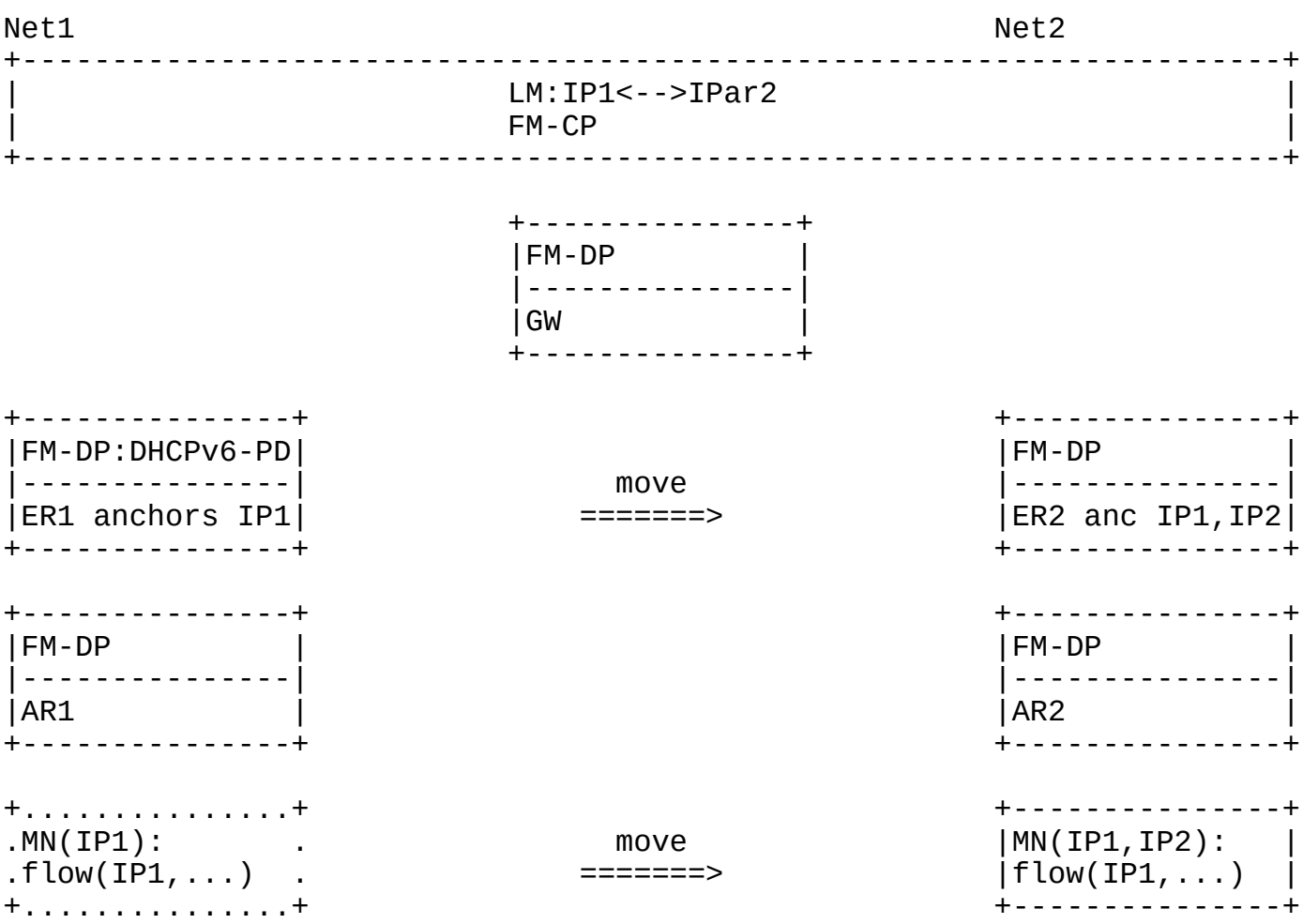

Figure 9. Mobility involving change of IP anchoring in a network with hierarchy in which the IP prefix allocated to the MN is anchored at an Edge Router supporting multiple access routers. to which the MN may connect to.

4. Security Considerations

TBD

5. IANA Considerations

This document presents no IANA considerations.

6. Contributors

This document is an attempt to harmonize the different distributed mobility solutions in a number of other drafts. These drafts cited in this document are the work of their many authors/co-authors. While some of them have taken the work to jointly write this

document, others have contributed at least indirectly by writing these drafts. The latter include Carlos J. Bernardos, Philippe Bertin, Hui Deng, Fabio Giust, Dapeng Liu, Satoru Matushima, Peter McCann, Antonio de la Oliva, Behcet Sarikaya, Pierrick Seite, Li Xue, Ryuji Wakikawa, and Younghan Kim.

Valuable comments have also been received from John Kaippallimil and ChunShan Xiong.

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